

## Stable homology-based cycle centrality measures



John Rick D. Manzanares\*, Paul Samuel P. Ignacio

*Department of Mathematics and Computer Science, University of the Philippines Baguio, Baguio City, 2600, Philippines.*

### Abstract

Network centrality measures play a crucial role in understanding graph structures, assessing the importance of nodes, paths, or cycles based on directed or reciprocal interactions encoded by vertices and edges. Estrada and Ross extended these measures to simplicial complexes to account for higher-order connections. In this work, we introduce novel centrality measures by leveraging algebraically-computable topological signatures of cycles and their homological persistence. We apply tools from algebraic topology to extract multiscale signatures within cycle spaces of weighted graphs, tracking homology generators persisting across a weight-induced filtration of simplicial complexes built over point clouds. This approach incorporates persistent signatures and merge information of homology classes along the filtration, quantifying cycle importance not only by geometric and topological significance but also by homological influence on other cycles. We demonstrate the stability of these measures under small perturbations using an appropriate metric to ensure robustness in practical applications. Finally, we apply these measures to fractal-like point clouds, revealing their capability to detect information consistent with, and possibly overlooked by, common topological summaries.

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### 1. Introduction

Many complex networks, such as social networks [10] and telecommunication networks [13] use centrality measures to determine the relative significance of nodes or cycles in the network. The derivations of the measures of central tendency in statistics reflect the idea that a single value can represent the entire distribution of a data set. In particular, the mode of data sets is comparable to the degree centrality of graphs in that it projects importance through frequencies. Similarly, closeness centrality is akin to the median of data in the sense that it identifies nodes that are reachable via short paths from any node as "central" nodes.

Giscard and Wilson [12] introduced the loop-centrality measure that uses the number of walks that intersect a loop to measure its importance. They found that this centrality measure has the ability to detect essential proteins in protein-protein interaction networks. For the same purpose, Estrada and Ross [9] explored an extension of this centrality measure to finite-dimensional substructures where higher-order connectivity networks (such as co-author [4] and social contact [14] networks) are represented by

\*Corresponding author

Email address: [jdmazanar@up.edu.ph](mailto:jdmazanar@up.edu.ph) (John Rick D. Manzanares)

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triangles and tetrahedra. A widely accepted notion of importance for cycles in a simplicial complex regards long-lived cycles as essential features of data, while short-lived cycles that appear are likely to be due to noise or sampling errors. However, Bubenik et al. [2] demonstrated that short-lived cycles hold important information that can be used to estimate the curvature of surfaces.

In this study, we propose novel centrality measures that leverage the persistence of homology classes and their merge history along the filtration. Integral to this is the development of an algorithm that captures the merge history of homology classes. These homology-based centrality measures produce, for all cycle generators, curves that reflect the relative importance of the corresponding generator throughout its entire evolution. By applying these centrality measures on toy models, we demonstrate the consistency of detected information by these measures to other topological summaries, and highlight its ability to capture new information possibly missed by other summaries. Accordingly, we establish some properties that include the stability of these measures under a distance analogous to norms in Lebesgue spaces and persistence landscapes.

## 2. Preliminaries

This section lays the groundwork for extracting multiscale features from weighted graphs. First, we model higher-order interactions among vertices in the graph using simplicial complexes, similar to the approach used in loop centrality [9]. These interactions capture relationships between vertices beyond simple pairwise connections. Next, we refine this model by applying a weight-based filtration. Filtration progressively removes less important edges, resulting in a nested sequence of increasingly coarser simplicial complexes. Finally, we leverage the concept of simplicial homology to analyze these filtered complexes. Simplicial homology allows us to identify and track cycle generators, which are loops within the complex that cannot be continuously deformed into a point. By focusing on these generators at each stage of the filtration, we can build a multiscale record of the graph's overall shape. This is crucial to designing our cycle centrality measures later. Throughout this chapter, we shall use [8, 11] as standard references for discussions involving simplicial and persistent homology.

### 2.1. Simplicial complex

We dive into the concept of simplicial homology and its application in modeling graphs. Although graphs represent connections between pairs of nodes (edges), simplicial complexes offer a more versatile framework. They allow us to capture not only pairwise interactions but also higher-order relationships between multiple nodes using simplices of various dimensions. This ability to encode complex interactions makes them particularly suitable for analyzing graphs with intricate connections beyond simple edges.

**Definition 2.1.** An (abstract) simplicial complex is a collection  $\mathcal{C}$  of subsets of a finite set  $V$  such that  $\tau \subseteq \sigma \in \mathcal{C}$  implies  $\tau \in \mathcal{C}$ . An element  $\sigma \in \mathcal{C}$  is called an (abstract) simplex with dimension  $|\sigma| - 1$ , and the largest such dimension among all simplices in  $\mathcal{C}$  defines the dimension of  $\mathcal{C}$ .

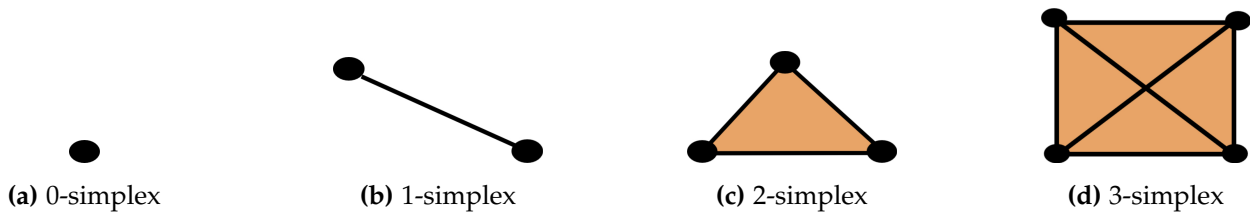


Figure 1: Lower-dimensional simplices in the Euclidean space.

In the Euclidean space, lower-dimensional simplices are named as *vertex*, *edge*, *triangle*, and *tetrahedron* for a 0-simplex, a 1-simplex, a 2-simplex, and a 3-simplex, respectively. Furthermore, the higher-dimensional simplices are polytopes analogous to triangles and tetrahedra.

When constructing a simplicial complex from a metric space, a common and less computational approach is through the *Vietoris-Rips complex* (or simply *Rips complex*). For a given set of points in a metric space and a fixed distance threshold, the Vietoris-Rips complex is the simplicial complex where any subset of points forms a simplex if all pairwise distances between points in the subset are below the threshold.

**Definition 2.2.** Let  $\epsilon > 0$ . Suppose that  $(M, d)$  is a metric space and  $P = \{p_i\}$  is a finite subset of  $M$ . The *Vietoris-Rips complex*  $\mathcal{R}_\epsilon(P)$  of  $P$ , with threshold  $\epsilon$ , is a simplicial complex whose  $k$ -simplices correspond to  $(k+1)$ -tuples of points in  $P$  such that  $d(p_i, p_j) \leq 2\epsilon$  for any pair of integers  $(i, j)$ .

Figure 2 illustrates a Rips complex of a point cloud with six initial points from the Euclidean space.

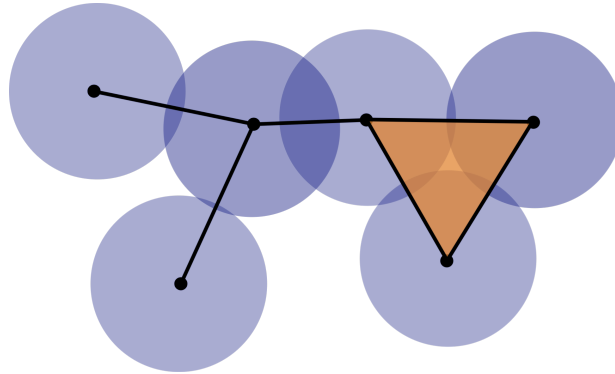


Figure 2: Rips complex of a point cloud.

The simplices of an abstract simplicial complex provide the fundamental building blocks for constructing chains in chain space. Each  $k$ -simplex acts as a single unit within a  $k$ -chain. By combining these simplices with coefficients from  $\mathbb{Z}/2\mathbb{Z}$  (which essentially encode the presence or absence of edges), we can create formal sums that represent more intricate relationships within the complex.

**Definition 2.3.** A  $k$ -chain is a formal sum of  $k$ -simplices with coefficients coming from the field  $\mathbb{Z}/2\mathbb{Z}$ . The *chain space*, denoted by  $C_k$ , is the free Abelian group generated by all possible  $k$ -simplices in the complex.

The field  $\mathbb{Z}/2\mathbb{Z}$  is selected to ignore the orientations of simplices, which simplifies calculations by treating oriented simplices as equivalent up to reordering. For example, an edge between vertices  $x_1$  and  $x_2$  is considered identical to the edge between  $x_2$  and  $x_1$ , as the orientation is disregarded.

Note that any  $k$ -simplex is a  $k$ -chain itself. For a further illustration of a  $k$ -chain, consider the simplicial complex in Figure 3. First, we denote a  $k$ -simplex, where  $k > 1$ , by  $[x_0, \dots, x_k]$  where  $\{x_0, \dots, x_k\}$  are collection of 0-simplices contained in the  $k$ -simplex. The formal sums  $[a, b] + [c, d]$  and  $[c, e] + [d, e]$  are examples of 1-chains, as all terms are 1-simplices. However, the expression  $[b, d] + [c, d, e]$  is not a  $k$ -chain because the terms belong to different dimensions:  $[b, d]$  is a 1-simplex, while  $[c, d, e]$  is a 2-simplex.

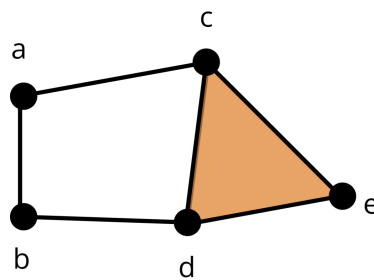


Figure 3: A simplicial complex (in the Euclidean space) with 0-simplices  $a$ ,  $b$ ,  $c$ ,  $d$ , and  $e$ .

These  $k$ -chains serve as the building blocks for studying the topological properties of the simplicial complex. However, not all  $k$ -chains represent actual cycles within the complex. To distinguish cycles from other chains, we introduce the concept of a boundary operator.

**Definition 2.4.** The *boundary operator*  $\partial_k : C_k \rightarrow C_{k-1}$  is the linear extension defined on the generators (individual  $k$ -simplices) given by

$$\partial_k ([x_0, \dots, x_j, \dots, x_k]) = \sum_{i=0}^k [x_0, \dots, \widehat{x_i}, \dots, x_k],$$

where  $\widehat{x_i}$  indicates that  $x_i$  is omitted.

Mathematically, the boundary operator acts on a  $k$ -simplex by summing over all its  $(k-1)$ -dimensional faces. For each face, it creates a new  $(k-1)$ -dimensional chain with the opposite coefficient. Intuitively, this captures the idea that the boundary of a  $k$ -dimensional object is made up of its  $(k-1)$ -dimensional faces. For instance, in the simplicial complex shown in Figure 3, we compute the boundary of a 1-chain as follows:

$$\partial_1 ([a, b] + [a, c]) = \partial_1 ([a, b]) + \partial_1 ([a, c]) = (b + a) + (c + a) = b + c.$$

Furthermore, observe that  $\partial_1 ([a, b] + [a, c] + [b, d] + [c, d]) = 0$ . Chains whose boundary maps to zero are called *cycles*, introduced in the following definition.

**Definition 2.5.** The *cycle space*  $Z_k$  and *boundary space*  $B_k$  of  $C_k$  are defined by  $Z_k = \ker \partial_k$  and  $B_k = \text{Im } \partial_{k+1}$ . We refer to elements of  $Z_k$  as  $k$ -cycles and those of  $B_k$  as  $k$ -boundaries.

Intuitively, these chains in the cycle space represent closed loops ( $k$ -cycles) within the complex that cannot be continuously deformed into a single point. On the other hand, chains in the boundary space correspond to the boundaries of higher-dimensional simplices.

The boundary operator, when applied twice, nullifies any  $k$ -boundary  $b_k$ . Formally, we have  $\partial_k \circ \partial_{k+1}(b_k) = 0$ . This property is known as the Fundamental Lemma of Homology [8, p. 95]. This lemma implies that every element of the boundary space is also an element of the cycle space, that is,  $B_k \subset Z_k$ . To illustrate this, consider the 2-simplex  $[c, d, e]$  in Figure 3. Its boundary is given by

$$\partial_2([c, d, e]) = [d, e] + [c, e] + [c, d],$$

and applying the boundary operator again yields  $\partial_1([d, e] + [c, e] + [c, d]) = 0$ .

## 2.2. Simplicial homology

Now that we can distinguish cycles and boundaries, we can introduce the concept of homology, which focuses on the essence of cycles, capturing their topological properties rather than their specific form.

**Definition 2.6.** Let  $\mathcal{C}$  be a simplicial complex. The  $k^{\text{th}}$  *homology group* of  $\mathcal{C}$ , written as  $H_k(\mathcal{C})$ , is the quotient group  $Z_k/B_k$ . Two cycles in  $Z_k(\mathcal{C})$  lying in the same homology class in this space are said to be *homologous*. We refer to the rank of  $H_k(\mathcal{C})$  as the  $k^{\text{th}}$  *Betti number* of  $\mathcal{C}$ .

In other words, the homology group is formed by taking the cycle space and "glue" any cycles that differ only by a boundary. This captures the intrinsic topological features of the complex at dimension  $k$ , independent of the specific choices of representatives for these cycles. The Betti number provides a numerical measure of the number of independent holes of dimension  $k$  in the complex.

By taking the quotient of the cycle space over the boundary space, we effectively exclude cycles that are simply the boundaries of higher-dimensional simplices. This ensures that we focus on nontrivial cycles that capture the true topological features of the complex. The result is a set of equivalence classes of cycles, with each class represented by a generator. These generators form a basis for the homology group,

encapsulating the independent  $k$ -dimensional features (loops or voids) of the space. Lastly, homologous cycles represent  $k$ -dimensional loops within the complex that can be continuously deformed into each other.

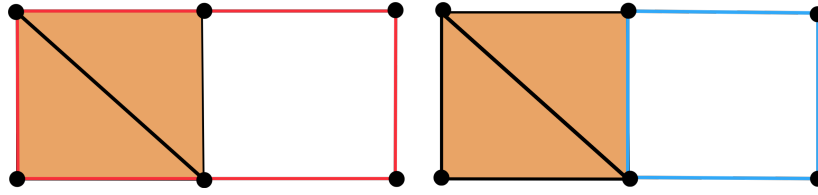


Figure 4: Two (red and blue) homologous 1-cycles.

The latter remark can be further visualized as follows. Consider two 1-cycles colored red and blue in the simplicial complex shown in Figure 4. These cycles may appear different, but if their difference can be expressed as the boundary of a 2-simplex in the complex, they are considered homologous. This implies that these cycles can be continuously deformed into each other, and any such cycle may serve as a representative of the homology class. The red and blue cycles, being homologous, are both potential generators for the first homology group of the complex, with one of them chosen to represent the essential 1-dimensional hole.

The generators of the homology group play a crucial role. They represent the distinct topological cycles embedded within our combinatorial model of the graph. Collectively, these cycles characterize the overall topology of the graph. Intriguingly, it is these generators and their corresponding homology classes that we leverage to define our centrality measures. However, a key point to remember is that equivalence within a homology class allows any generator to represent its class. For example, in Figure 4, either the red or blue 1-cycle can represent the same topological hole. This raises a natural question: which cycle becomes the representative of the homology class? This is where the concept of persistence comes into play, which we will explore in the next section. For succeeding sections, we use the notation  $[\sigma]$  for a homology class represented by the cycle  $\sigma$ .

### 2.3. Filtrations and Persistent Homology

This section explores filtrations and persistent homology, powerful tools for analyzing topological features within simplicial complexes. These tools allow us to track how the underlying topological structure of a complex evolves as we progressively remove or modify its simplices based on some criteria.

We begin by introducing the concept of weights in simplicial complexes. A weight function  $w : \mathcal{C} \rightarrow \mathbb{R}_{\geq 0}$  assigns weights (nonnegative real numbers) to each simplex  $\sigma$  in the complex  $\mathcal{C}$ . This function allows us to differentiate between simplices and prioritize their removal during the filtration process. A common example of a weight function is the distance function, where the weight of a simplex might be defined as the maximum distance between any two points within that simplex. Higher weight values would then correspond to simplices with larger maximum distances, potentially representing less connected or more spread-out regions of the complex.

A filtration is a nested sequence of subcomplexes obtained by progressively removing simplices exceeding a certain weight threshold. Different weight order (e.g.  $\geq$  order) leads to different filtration types. Intuitively, as the weight threshold increases, simplices with lower importance (based on the weight function) are removed. This reveals a coarser version of the original complex at each step, where only the most significant simplices remain.

This filtration process induces a sequence  $\{\mathcal{C}_{w_i}\}$  of simplicial complexes at thresholds  $w_i$ . These complexes capture the connectivity structure of the original complex at each weight threshold. Persistent homology leverages this sequence to track the "persistence" of topological features (cycles) across different weight thresholds. In simpler terms, we are interested in how these features are born, die, or evolve as the filtration progresses.

**Definition 2.7.** Let  $\{\mathcal{C}_{w_i}\}$  be a filtration of a simplicial complex  $\mathcal{C}$  at thresholds  $w_i$ . The  $k^{\text{th}}$   $w_i - w_j$  persistent homology group of  $\mathcal{C}$ , denoted by  $H_k^{w_i, w_j}$ , is the image of the induced homomorphism

$$H_k(\mathcal{C}_{w_i}) \xrightarrow{f_k^{w_i, w_j}} H_k(\mathcal{C}_{w_j}).$$

This group captures the homology classes "born" at weight  $w_i$  and persisting (remaining in the homology group) until weight  $w_j$ . This group encodes how long homology classes survive under the filtration. For instance, a cycle that persists across a wide range of weight values suggests a more robust topological feature in the complex.

For the following definitions, we let  $w_i, w_j$ , and  $w_k$  be real numbers such that  $w_i < w_j < w_k$ .

**Definition 2.8.** Let  $[\sigma]$  be a homology class. We say that  $[\sigma]$  is *born* at  $\mathcal{C}_{w_j}$  if  $[\sigma] \in H_k^{w_j, w_k}$  but  $[\sigma] \notin H_k^{w_i, w_j}$ . The *birth*  $b(\sigma)$  of  $[\sigma]$  is given by  $w_j$ . Furthermore, if  $[\sigma]$  is born at  $\mathcal{C}_{w_i}$ , then we say that  $[\sigma]$  *dies* entering  $\mathcal{C}_{w_k}$  if  $f_k^{w_i, w_k}([\sigma]) = 0$  but  $f_k^{w_i, w_j}([\sigma]) \neq 0$ . The *death*  $d(\sigma)$  of  $[\sigma]$  is given by  $w_k$ .

A homology class dies when it becomes identical to a combination of classes that are born earlier. In this context, the elder rule comes into play. This rule selects the generator formed at the lower threshold as the natural representative of the merged class. A key observation is made: even when a different generator survives, persistence can be "transferred" to the surviving one. This highlights the dynamics of merging classes and motivates further investigation into how these classes evolve under the filtration. The notion of  $k$ -nearness formalizes the idea that two cycles are close in the complex.

**Definition 2.9.** Let  $\sigma$  and  $\nu$  be  $k$ -cycles. We say that  $\sigma$  and  $\nu$  are  $k$ -near if there exists a  $k$ -chain common to both  $\sigma$  and  $\nu$ . Furthermore, the  $k^{\text{th}}$  homology classes  $[\sigma]$  and  $[\nu]$  are  $k$ -near if there exist  $\sigma \in [\sigma]$  and  $\nu \in [\nu]$  that are  $k$ -near.

We also define the concept of merging governed by the elder rule. This definition clarifies the direction of merging, where a class with an earlier death threshold merges with the long-lived class.

**Definition 2.10.** Let  $[\sigma]$  and  $[\nu]$  be distinct  $k^{\text{th}}$  homology classes in a simplicial complex  $\mathcal{C}_\epsilon$ , where  $d(\sigma) \neq d(\nu)$ . We say that  $[\sigma]$  and  $[\nu]$  *merge* at time  $\epsilon' = \min\{d(\sigma), d(\nu)\}$  if

$$\sigma + \nu = \rho \tag{2.1}$$

for some  $k$ -boundary  $\rho$ .

Note that Definition 2.10 is well-defined in the sense that Equation (2.1) still holds regardless of the choice of representatives for the homology classes. To demonstrate this, let  $\sigma' \in [\sigma]$  be any representative of the homology class  $[\sigma]$ . Then,  $\sigma$  can be expressed as  $\sigma = \sigma' + \rho$ , where  $\rho$  is some  $k$ -boundary. Similarly, for any representative  $\nu' \in [\nu]$ , we have  $\nu = \nu' + \rho'$  for some  $k$ -boundary  $\rho'$ . If  $\sigma + \nu$  is a  $k$ -boundary, say  $\rho''$ , then it follows that

$$\sigma' + \rho + \nu' + \rho' = \rho''.$$

This implies that  $\sigma' + \nu' = \rho_0$ , where  $\rho_0 = \rho + \rho' + \rho''$  is a  $k$ -boundary. Thus, the choice of representatives does not affect the outcome, establishing the well-definedness of Definition 2.10.

It is crucial to distinguish between the merging process and the act of gluing a boundary to a cycle. While merging signifies the disappearance of a topological feature upon simplex removal, gluing a boundary to a cycle essentially fills in the hole or closes the gap within the cycle. Although the resulting structure remains homologous, it falls outside the scope of current persistent homology methods. Addressing this distinction involves investigating the sequential gluing actions across the filtration.

The dynamic nature of merging classes is underscored by the observed phenomenon of information transfer and the need to resolve boundary gluing. These observations prompt further inquiry into the evolution of these classes under the filtration process.



#### 2.4. Algorithm for class merging

The computation of persistent homology often involves the matrix reduction algorithm introduced by Zomorodian [19]. This algorithm identifies the representatives for each homology class by reducing the boundary matrices of the simplicial complex. The  $k^{\text{th}}$  boundary matrix  $\partial_k$  is constructed by arranging the  $(k-1)$ -simplices as rows and the  $k$ -simplices as columns, sorted in increasing order of birth in the filtration. The entries of the matrix are determined using the map

$$\partial_k[i, j] = \begin{cases} 1, & \text{if } \sigma_i \text{ is a codimension-1 face of } \sigma_j, \\ 0, & \text{otherwise.} \end{cases}$$

Each column  $j$  in the boundary matrix corresponds to a  $k$ -simplex  $\sigma_j$ , and each row  $i$  corresponds to a  $(k-1)$ -simplex  $\sigma_i$ .

A pseudocode for the matrix reduction algorithm [8, p. 104] is presented in Algorithm 1. The key operation in this algorithm is to ensure that each column has a unique pivot, which is the lowest non-zero entry in the column. The notation  $\text{low}(k)$  refers to the row index of the lowest 1 in column  $k$  of the boundary matrix.

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**Algorithm 1** Standard reduction algorithm for boundary matrices.

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**Require:** Boundary matrix of order  $m$

**Ensure:** Reduced boundary matrix

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1: for  $i = 2$  to  $m$  do
2:   while there exists  $i < j$  with  $\text{low}(i) = \text{low}(j)$  do
3:     add column  $i$  to column  $j$ 
4:   end while
5: end for
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Let  $\text{col}(\sigma)$  denote the column vector associated with the simplex  $\sigma$  in the boundary matrix. After the matrix is reduced, each column is expressed as a chain, and we denote the chain corresponding to  $\sigma$  as  $\text{rep}(\sigma)$ . Columns that reduced to a zero vector represent homology classes.

For the proof of the following lemma, we adopt an abuse of notation by defining the appearance or birth of a simplex  $\sigma$  by  $b(\sigma)$ . In this context, we say that a simplex  $\sigma_i$  *gives birth to* the cycle  $\sigma$  if  $b(\sigma) = b(\sigma_i)$ . Additionally, we write  $\sigma \prec \delta$  to indicate that either  $b(\sigma) < b(\delta)$ , or, if  $b(\sigma) = b(\delta)$ ,  $\text{col}(\sigma)$  appears before  $\text{col}(\delta)$  in the boundary matrix.

**Lemma 2.11.** *Suppose that  $[\sigma]$  is a  $k^{\text{th}}$  homology class in a filtered simplicial complex  $\mathcal{C}$ . If  $[\sigma]$  is not  $k$ -near to any homology class  $[\nu]$ , then  $[\sigma]$  has a unique class representative produced by Algorithm 1.*

*Proof.* We assume that  $\sigma = \sum_i \sigma_i$  and  $\nu = \sum_j \nu_j$  are not  $k$ -near cycles where some  $\sigma_i$  and  $\nu_j$ , respectively, give birth to  $\sigma$  and  $\nu$ . Furthermore, let  $\nu_j \prec \sigma_i$ . Note that the  $k^{\text{th}}$  homology class representative depends on the associated  $k^{\text{th}}$  boundary matrix, which is only concerned with simplices of dimensions  $k-1$  and  $k$ . Therefore, we can omit the case where  $\sigma$  and  $\nu$  has a common simplex of dimension 0 to  $k-2$ .

Suppose that  $\sigma$  and  $\nu$  have no shared  $(k-1)$ -simplex. Then, for every simplex  $\delta$ ,  $\text{low}(\delta) \neq \text{low}(\sigma_i)$  for every integer  $1 \leq i \leq n$ . Therefore, the representative of  $[\sigma]$  must be unique, since no simplex can be added with  $\sigma_i$  to form another representative.

Now, consider the case where  $\sigma$  and  $\nu$  intersect at some  $(k-1)$ -faces. For the sake of contradiction, suppose that  $\sigma$  has a representative containing some simplices from  $\nu$ . Since  $\sigma$  and  $\nu$  has no shared  $k$ -simplex,  $\text{rep}(\sigma)$  must contain all simplices of  $\nu$ . Otherwise,  $\text{rep}(\sigma)$  ceases to be a cycle if it does not contain all simplices of  $\nu$ . Note that, in the reduction algorithm,  $\text{col}(\nu_j)$  reduces to a zero vector before reduction happens on  $\text{col}(\sigma_i)$ . Hence,  $\text{col}(\sigma_i)$  cannot contain the simplex  $\nu_j$  of  $\nu$ . Therefore, we arrive at a contradiction so  $[\sigma]$  must have a unique class representative.  $\square$

To illustrate Lemma 2.11, consider the simplicial complex shown in Figure 5b. The homology class  $[\sigma]$ , where  $\sigma = [c, d] + [c, e] + [d, e]$ , is not  $k$ -near to any other homology class. This implies that  $\sigma$  is the unique representative for this homology class. This is evident by following the reduction algorithm, even when edges connect to the hole without forming another cycle. For instance, assume the ordering  $a \prec b \prec d \prec e \prec c$  and  $[a, c] \prec [c, d] \prec [c, e] \prec [d, e]$ . Initially,  $\text{rep}([a, c])$  is added to  $\text{rep}([c, d])$  and  $\text{rep}([c, e])$ . Then these two columns are added to  $\text{rep}([d, e])$ , resulting in

$$\text{rep}([d, e]) = [d, e] + ([a, c] + [c, e]) + ([a, c] + [c, d]) = [d, e] + [c, e] + [c, d].$$

If  $[\sigma]$  were 1-near to another hole, there would be multiple potential representatives for the homology class produced by the reduction algorithm. Further details are discussed in Lemma 2.12.

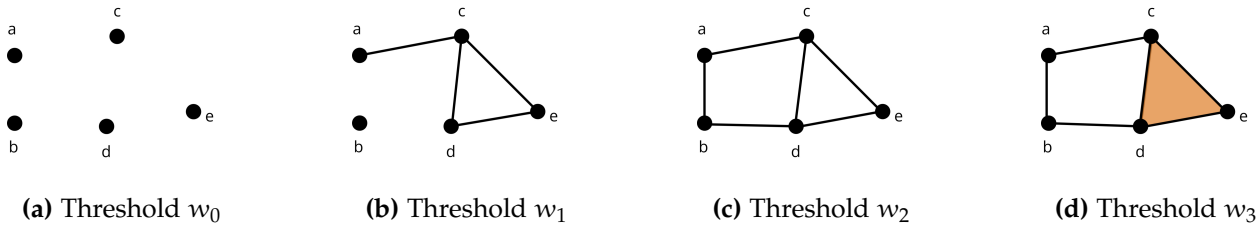


Figure 5: Filtration of a simplicial complex at multiple weight thresholds  $w_0, w_1, w_2$ , and  $w_3$  satisfying  $w_0 < w_1 < w_2 < w_3$ .

Consider two  $k$ -near cycles  $\sigma$  and  $\nu$  that share a chain  $\delta$ . The shared simplices, along with simplices from  $\sigma$  and  $\nu$  that touch these shared simplices, can create situations where some simplices in  $\sigma + \delta$  might have the same position of the lowest one as some simplices in  $\nu + \delta$  and  $\delta$ . We call such connection points as *junctures*.

**Lemma 2.12.** *Suppose that  $[\sigma]$  and  $[\nu]$  are distinct  $k^{\text{th}}$  homology classes appearing in the filtration of a simplicial complex satisfying  $\nu \prec \sigma$ . Let  $\nu = \sum_p \nu_p$  be a cycle representing  $[\nu]$ , where  $\nu_p$  gives birth to  $\nu$ . Additionally, assume  $\sigma$  and  $\nu$  are  $k$ -near with common chain  $\delta$ . If  $\nu_p \in \delta$  and  $\nu$  does not contain simplices of  $\sigma$  except  $\delta$ , then  $\text{rep}(\sigma)$  is  $\sigma + \nu$ . Otherwise, if  $\nu_p \notin \delta$ , then  $\text{rep}(\sigma)$  is  $\sigma$ .*

*Proof.* Suppose that  $\sigma = \sum_i \sigma_i$  and  $\nu$  are  $k$ -near with the intersection  $\delta$ , and  $\nu_p \in \delta$ . Let  $\sigma_i$  be the simplex giving birth to  $\sigma$ . Note that  $\text{rep}(\sigma_i)$  only depends on the junctures and the algorithm proceeds by having  $\sigma + \delta$  in  $\text{rep}(\sigma_i)$ . If  $\delta \in \text{rep}(\sigma)$ , then we arrive at a contradiction since  $\nu_p \in \delta$  and  $\nu_p$  gives birth to  $\nu$  so  $\text{col}(\nu_p)$  is a zero vector.

If there exists no juncture where  $\sigma_j \in \sigma + \delta$  and  $\nu_q \in \nu$  such that  $\text{low}(\sigma_j) = \text{low}(\nu_q)$ , then  $\text{rep}(\sigma_i)$  ceases to be a cycle. Thus, there exists  $\sigma_j \in \sigma + \delta$  and  $\nu_q \in \nu$  such that  $\text{low}(\sigma_j) = \text{low}(\nu_q)$ . By the algorithm,  $\text{rep}(\nu_q) \in \text{rep}(\sigma_i)$ , and we have  $\nu + \delta \in \text{rep}(\sigma_i)$ . Consequently,

$$\text{rep}(\sigma_i) = (\sigma + \delta) + (\nu + \delta) = \sigma + \nu.$$

Therefore, we have proved the lemma.  $\square$

In the simplicial complex shown in Figure 5c, there are 1-near classes  $[\sigma]$  and  $[\delta]$ , where  $\sigma = [a, c] + [c, d] + [a, b] + [b, d]$  and  $\delta = [c, d] + [c, e] + [d, e]$ . Two possible candidates exist for the representative of the hole  $[\sigma]$ : either  $\sigma$  itself or  $\sigma + \delta$ . According to Lemma 2.12,  $\sigma$  is selected as the representative by Algorithm 1 if  $[c, d] \prec [c, e], [d, e]$ . However, if  $[c, e], [d, e] \prec [c, d]$ , then the representative becomes  $\sigma + \delta$ . In short, when  $[c, e], [d, e] \prec [c, d]$ , the column corresponding to  $[c, d]$  reduces to a zero vector, meaning  $[c, d]$  is not added to any other columns on its right due to the absence of a lowest one position.

Algebraically, from Lemma 2.12, we can obtain  $\sigma + \nu$  from  $\nu$  using a change of basis technique. This operation corresponds to adding two zero column vectors while updating  $\text{rep}(\sigma)$ . When applying this technique specifically to  $k$ -near classes, we encounter a concept analogous to short and long path cycles in graph theory. In this setting, the algorithm allows us to select a class representative based on the computed representatives.



**Theorem 2.13.** *Let  $[\sigma]$  and  $[\nu]$  be distinct  $k^{\text{th}}$  homology classes, where  $\sigma \prec \nu$ . Then  $[\sigma]$  merges with  $[\nu]$  at  $\epsilon = d(\sigma)$  if and only if  $[\sigma]$  and  $[\nu]$  are  $k$ -near.*

*Proof.* Suppose  $[\sigma]$  merges with  $[\nu]$ . If  $[\sigma]$  and  $[\nu]$  are not  $k$ -near, then both classes have unique cycle representatives that are not also  $k$ -near. Thus, if  $\sigma + \nu$  is a boundary at  $\epsilon$ , then  $d(\sigma) = d(\nu)$ . This equation is a contradiction to the assumption that the classes merge. Therefore,  $[\nu]$  and  $[\sigma]$  must be near.

Conversely, suppose that  $[\sigma]$  and  $[\nu]$  are  $k$ -near. We can respectively choose  $\sigma + \nu$  and  $\nu$  as representatives of  $[\sigma]$  and  $[\nu]$ . Observe that  $(\sigma + \nu) + \nu = \sigma$ . Since  $\sigma \prec \nu$ ,  $\sigma$  must be a boundary at  $\epsilon$ . Therefore,  $[\sigma]$  merges with  $[\nu]$ .  $\square$

According to Theorem 2.13, the two holes in Figure 5c will merge as they share the edge  $[c, d]$ . This statement also holds using Definition 2.10. Observe that, we can represent the hole formed by the 0-simplices  $a, b, c$ , and  $d$  as  $[a, c] + [a, b] + [b, d] + [c, e] + [d, e]$  using a change of basis technique on the representatives. After the death of the hole  $[c, d] + [c, e] + [d, e]$ , the two holes become homologous, as their sum equals  $\partial_2([c, d, e])$ .

In the filtration illustrated in Figure 5, consider the scenario where the cycle  $cf + fg + eg + ce$  is born at a threshold exceeding  $w_3$ . By applying a change of basis to the existing cycle (given birth by  $[b, d]$ ) with vertices  $a, b, c$ , and  $d$ , we can obtain the cycle  $ab + ac + bd + de + cf + fg + eg$  as its representative. Suppose that the edge  $[a, d]$  appears, followed by the triangles  $[a, c, d]$  and  $[a, b, d]$ . According to the reduced boundary matrix for this filtration,  $\text{col}([a, b, d])$  has a lowest one at row corresponding to edge  $[b, d]$ . This means that the appearance of triangle  $[a, b, d]$  destroys the hole represented by  $\text{rep}([b, d])$ . Hence, the newly introduced cycle persists while the elder hole vanishes. Although this approach diverges from the elder rule, the cycles become homologous at their respective thresholds. Therefore, to facility this study, we propose a relaxation of the elder rule while maintaining the fundamental definition of merging within the persistent homology framework.

In addition to keeping track of the cycle representatives yielded by the reduction algorithm, we are also interested in integrating the *snowball effect* of successive merging classes in persistent homology. This necessitates the following definition.

**Definition 2.14.** Consider the induced homomorphism

$$H_k(\mathcal{C}_{w_i}) \xrightarrow{f_k^{w_i, w_j}} H_k(\mathcal{C}_{w_j})$$

and let  $[\sigma] \in H_k(\mathcal{C}_{w_j})$ . We define the *first order merge cluster* of  $[\sigma]$  at  $w_j$  as

$$M_1[\sigma, w_j] := (f_k^{w_i, w_j})^{-1}([\sigma]),$$

that is, the set of homology classes merging with  $[\sigma]$  at threshold  $w_j$ . Inductively, for every integer  $n \geq 2$ , the  $n^{\text{th}}$  *order merge cluster* of  $[\sigma]$  at  $w_j$  is defined as

$$M_n[\sigma, w_j] = \bigcup_{\tau \in M_{n-1}[\sigma, w_j]} M_1[\tau, w_j].$$

The first-order merge clusters refer to all other classes that merges with a specific class  $[\sigma]$  across the filtration until the threshold  $w_j$ . Higher-order merge clusters build on this idea. Imagine a class that merges with another class, which itself merges with a third class. All these classes are considered part of the same merge cluster because their merges are ultimately connected. The definition captures this cascading effect by recursively defining higher-order merge clusters as the union of all first-order merge clusters of classes within the previous-order merge cluster.

Theorem 2.13 offers a way to determine if two classes merge at a specific threshold. However, it might involve checking many pairs of cycles representing the classes, which can be computationally expensive. This is because any homologous cycle can represent a class, leading to a large number of potential pairings to explore. Corollary 2.15 aims to reduce the number of pairings that need to be checked for merging. It leverages the concept of  $k$ -nearness between cycles representing classes.

**Corollary 2.15.** *Suppose that the homology class  $[\sigma]$  merges with  $[\nu]$ . Let  $[\delta]$  be a homology class such that  $\sigma \prec \delta \prec \nu$ . If  $[\nu]$  and  $[\delta]$  are near then either  $[\delta]$  merges with  $[\sigma]$ , or there exists a homology class that merges with  $[\nu]$  whose  $n^{\text{th}}$  order merge cluster contains  $[\delta]$  for some integer  $n \geq 1$ .*

*Proof.* Suppose  $[\nu]$  and  $[\delta]$  are near. If  $[\delta]$  merges with  $[\nu]$ , then we have proved the corollary. Suppose  $[\delta]$  merges with a homology class  $[\lambda]$ , where  $\delta \prec \lambda \prec \nu$ . Note that a representative of  $[\lambda]$  includes some simplices of  $\delta$ . Since  $[\nu]$  and  $[\delta]$  are near,  $[\nu]$  and  $[\lambda]$  must also be near. Therefore,  $[\lambda]$  merges with  $[\sigma]$ .  $\square$

We can implement Theorem 2.13 and Corollary 2.15 in an algorithm to obtain the first-order merge clusters of each homology class. Algorithm 2 presents a pseudocode for finding the first order merge clusters.

---

**Algorithm 2** Identifying first-order merge clusters from class representatives.

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**Require:** Representatives  $\{\sigma_i\}_{i=1}^l$  ordered by ascending birth followed by death

**Ensure:** Collection containing first order merge clusters of the homology classes

```

1: Initialize an array M of length l
2: for  $i \in \{2, \dots, l\}$  do
3:   for  $j \in \{1, \dots, i-1\} \setminus \bigcup_{k=1}^{i-1} M[k]$  do
4:     if  $[\sigma_i]$  is k-near to  $[\sigma_j]$  and  $d(\sigma_i) \geq d(\sigma_j) > b(\sigma_i)$  then
5:       insert j to M[i]
6:     end if
7:   end for
8: end for

```

---

### 3. Cycle centrality

In this chapter, we propose centrality measures for holes based on persistent homology. These measures go beyond their mere existence and capture its topological significance and influence within an object. We achieve this by considering two key aspects, namely persistence and merge dynamics.

#### 3.1. Centrality measures

Three centrality functions, denoted by  $J_1$ ,  $J_2$ , and  $J_3$ , are defined to capture the evolving importance of cycles. Each function is defined for a specific homology class representative  $[\sigma]$  and a filtration threshold  $\epsilon$ . All the functions share a common structure. They possess a base value of 0 before the birth threshold of  $[\sigma]$ , and exhibit a piece-wise linear behavior after birth, reflecting changes in importance.

The first centrality measure  $J_1$  aims to capture and quantify the total persistence accumulated by  $[\sigma]$  and all classes that directly merge with it up until the threshold  $\epsilon$ . Hence, if  $P_\epsilon(\sigma)$  is the persistence of  $\sigma$  at  $\epsilon$ , the first centrality function has the form

$$J_1(\sigma, \epsilon) = \begin{cases} 0, & \text{for } \epsilon \leq b(\sigma), \\ P_\epsilon(\sigma) + \sum_{[\phi] \in M_1[\sigma, \epsilon]} P_\epsilon(\phi), & \text{for } \epsilon > b(\sigma). \end{cases}$$

As each  $P_\epsilon(\sigma)$  is monotonic and stabilizes when  $\epsilon > d(\sigma)$ , this function is piece-wise linear and monotonic. It captures the cascading simple aggregate of persistence pooled from cycle representatives that altogether merge directly to an older cycle. We view this as the homological importance of cycles -if many cycles merge to an old cycle, then its homological significance is proportionally increased.

One caveat of the function above is that it treats all merging instances similarly regardless of when the merge happens along the filtration. Hence, we also consider a second centrality function  $J_2$  by refining  $J_1$  to account for the time when instances of merging happen. It introduces a scaling function that assigns a

weight to the persistence of each merging class. This allows us to prioritize either early or late merges in the centrality calculation. We can write this function as

$$J_2(\sigma, \epsilon) = \begin{cases} 0, & \text{for } \epsilon \leq b(\sigma), \\ P_\epsilon(\sigma) + \sum_{[\phi] \in M_1[\sigma, \epsilon]} f_\sigma(\phi) P_\epsilon(\phi), & \text{for } \epsilon > b(\sigma), \end{cases}$$

where  $f_\sigma(\phi)$  is a scaling factor dependent on  $[\sigma]$  and  $[\phi]$ . When early merges are considered more influential, we can define  $f_\sigma(\phi)$  as  $d(\phi)/d(\sigma)$ . Conversely, when late merges are considered more important, we can define  $f_\sigma(\phi)$  as  $1 - d(\phi)/d(\sigma)$ .

We can also generalize the cascading effect of merging by considering the indirect transfer of persistence from merging instances prior to a given merge time. This is equivalent to modifying the centrality function to account for higher-order merging clusters, and yields a third centrality function  $J_3$  given by

$$J_3(\sigma, \epsilon) = \begin{cases} 0, & \text{for } \epsilon \leq b(\sigma), \\ P_\epsilon(\sigma) + \sum_r \sum_{[\phi] \in M_r[\sigma, \epsilon]} f_\sigma(\phi) P_\epsilon(\phi), & \text{for } \epsilon > b(\sigma). \end{cases}$$

Here, we allow the definition  $f_\sigma(\phi) = 1$  to generalize the function  $J_1$ .

In general, the centrality measures are of the form  $J_n : \Lambda \times W \rightarrow \mathbb{R}_+$ , where  $\Lambda$  is the collection of all homology classes from a filtration with a nonzero persistence. Hence, each persistence diagram produces a family of centrality functions. We can visualize the time-evolving centrality function of each homology class by constructing a piecewise-linear plot where sudden increases represent the effect of merging. We call these plots the  $J_n$  *centrality plots* of dimension  $k$ , where  $k$  refers to the dimension of the homology classes we are concerned with. We can also create heat maps from contiguous heat bars each capturing the monotonic growth of centrality as captured by our measures. We illustrate a centrality plot in Figure 6 and omit the heat map representation.

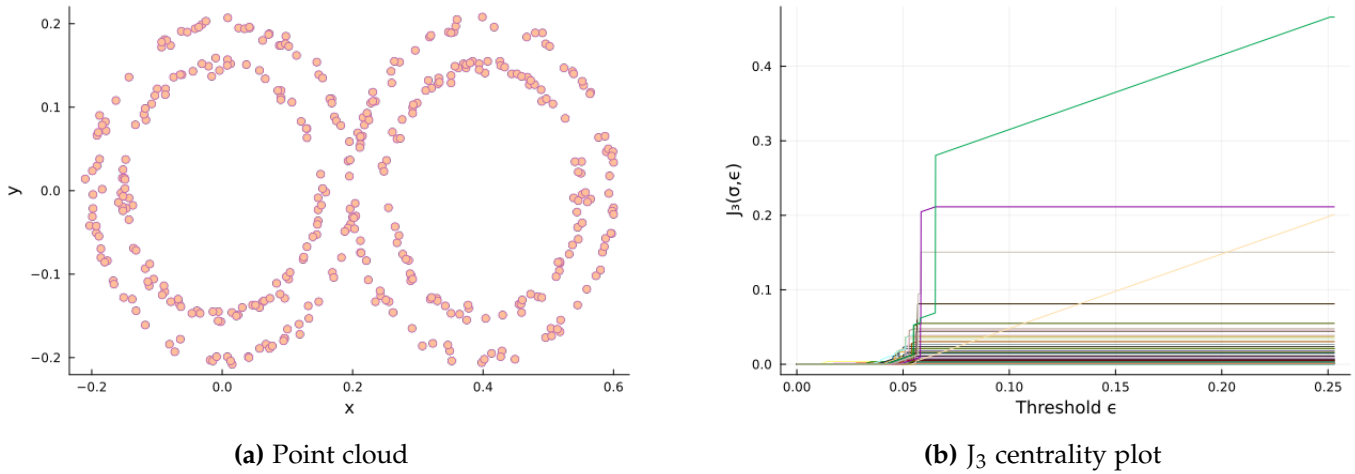


Figure 6: The  $J_3$  centrality plot, with  $f_\sigma = 1$ , of dimension 1 produced by the Rips filtration of the point cloud sampled around a wedge sum of two annuli.

### 3.2. Stability of centrality measures

This subsection investigates the stability of the centrality measures we defined earlier. Stability ensures that small changes in the network data, such as slight adjustments to edge weights, will not lead to drastic changes in the calculated centrality of cycles.

First, we establish that all three centrality measures are monotonic. Intuitively, as we explore the network at a coarser scale, a hole might encounter more merging partners, potentially accumulating more persistence and thus increasing its centrality. In the following lemma, we denote  $J_n(\sigma, \epsilon)$  by  $J_{n,\sigma}(\epsilon)$ .

**Lemma 3.1** (Monotonicity). *If  $\epsilon < \epsilon'$ , then  $J_{n,\sigma}(\epsilon) \leq J_{n,\sigma}(\epsilon')$ .*

*Proof.* It follows from definition that  $M_k[\sigma, \epsilon] \subseteq M_k[\sigma, \epsilon']$  for  $\epsilon < \epsilon'$  and  $k \geq 1$ . Since  $P_\epsilon$  is monotonic and  $f$  is constant with respect to  $\epsilon$ , the conclusion follows.  $\square$

To analyze stability, we introduce the concept of a  $p$ -centrality norm. This norm assigns a numerical value to each hole's centrality, capturing the overall importance of that hole. The specific value of  $p$  influences how the norm prioritizes different aspects of the centrality function. For example,  $p = 1$  emphasizes the average centrality, while  $p = \infty$  focuses on the maximum centrality value achieved by the hole.

**Definition 3.2.** Let  $\mathcal{J}_n = \{J_{n,\sigma} | [\sigma] \in \Lambda\}$  denote the collection of centrality functions generated by the set of persistent homology classes  $\Lambda$ . The  $p$ -centrality norm is given by

$$\|J_{n,\sigma}(d^*)\|_p = \begin{cases} \left( \int_0^{d^*} (J_{n,\sigma}(x))^p dx \right)^{1/p}, & \text{if } 1 \leq p < \infty, \\ J_{n,\sigma}(d^*), & \text{if } p = \infty, \end{cases}$$

where  $d^*$  is the minimum between  $d(\sigma)$  and the largest geodesic distance between any two vertices in the largest cycle in  $[\sigma]$ .

We next define a metric called the  $p$ -centrality distance, which is used to compare collections of centrality norms. Since a natural order does not exist for these centrality measures, we adopt a bottleneck-like distance approach. This method computes optimal pairings between centrality norms in two collections, focusing on the maximum difference between any two paired values. To enhance computational efficiency, we follow the condition outlined in the bottleneck distance implementation from [15], where the computation is simplified by assuming that the birth of each class is the same in both collections.

Formally, let  $\Omega = \{0\} \times \{\|J_{n,\sigma}(d^*)\|_p^p : J_{n,\sigma} \in \mathcal{J}_n\}$  and  $\Omega' = \{0\} \times \{\|J_{n,\sigma'}(d^{*'})\|_p^p : J_{n,\sigma'} \in \mathcal{J}'_n\}$ . For  $x_\sigma \in \Omega$ , we define  $\delta_{x_\sigma} = \|J_{n,\sigma}\|_p^p$ . Given a bijection  $\phi : \Omega \cup \Delta \rightarrow \Omega' \cup \Delta$ , where  $\Delta = \{(x, x) : x \in \mathbb{R}\}$ , the  $p$ -centrality distance is given by

$$\|x_\sigma - \phi(x_\sigma)\|_\infty = \begin{cases} \frac{1}{2} \max\{\delta_{x_\sigma}, \delta_{\phi(x_\sigma)}\}, & \text{if } \phi(x_\sigma) \in \Delta, \\ |\delta_x - \delta_{\phi(x)}|, & \text{otherwise.} \end{cases} \quad (3.1)$$

**Definition 3.3.** For  $1 \leq p < \infty$ , the  $p$ -centrality distance is given by

$$C_p(\mathcal{J}_n, \mathcal{J}'_n) = \inf_{\phi} \sup_{x \in X} \|x - \phi(x)\|_\infty,$$

where the infimum is taken over all bijections from  $\Omega \cup \Delta$  to  $\Omega' \cup \Delta$ .

For the case where  $p = \infty$ , we propose a distance akin to  $p$ -landscape distance [3, p. 94]. Note that we can order the centrality function  $J_{n,\sigma}(d^*)$  of each homology class  $[\sigma]$  based on the maximum centrality values. Thus, we obtain a decreasing sequence  $\{\|J_{n,m}\|_m\}$ , where  $m$  is a positive integer. The  $p$ -centrality distance is then given by

$$C_p(\mathcal{J}_{n,k}, \mathcal{J}'_{n,k}) = \sum_m \|J_{n,m} - J'_{n,m}\|_p.$$

**Example 3.4.** We examine how our proposed centrality measures behave with respect to perturbations of the point cloud in Figure 6a introduced by replacing each point  $(x, y)$  with  $(x + \kappa_1, y + \kappa_2)$  for some  $\kappa_1, \kappa_2 \in [-\kappa, \kappa]$ . We then compute the 1-centrality distance between the centrality measures of the original point cloud and its perturbation, and replicate thirty simulations of this process to generate a distribution of 1-centrality distances represented by boxplots. In Figures 7, we observe how the distribution of 1-centrality distances varies across increasing levels of perturbations.

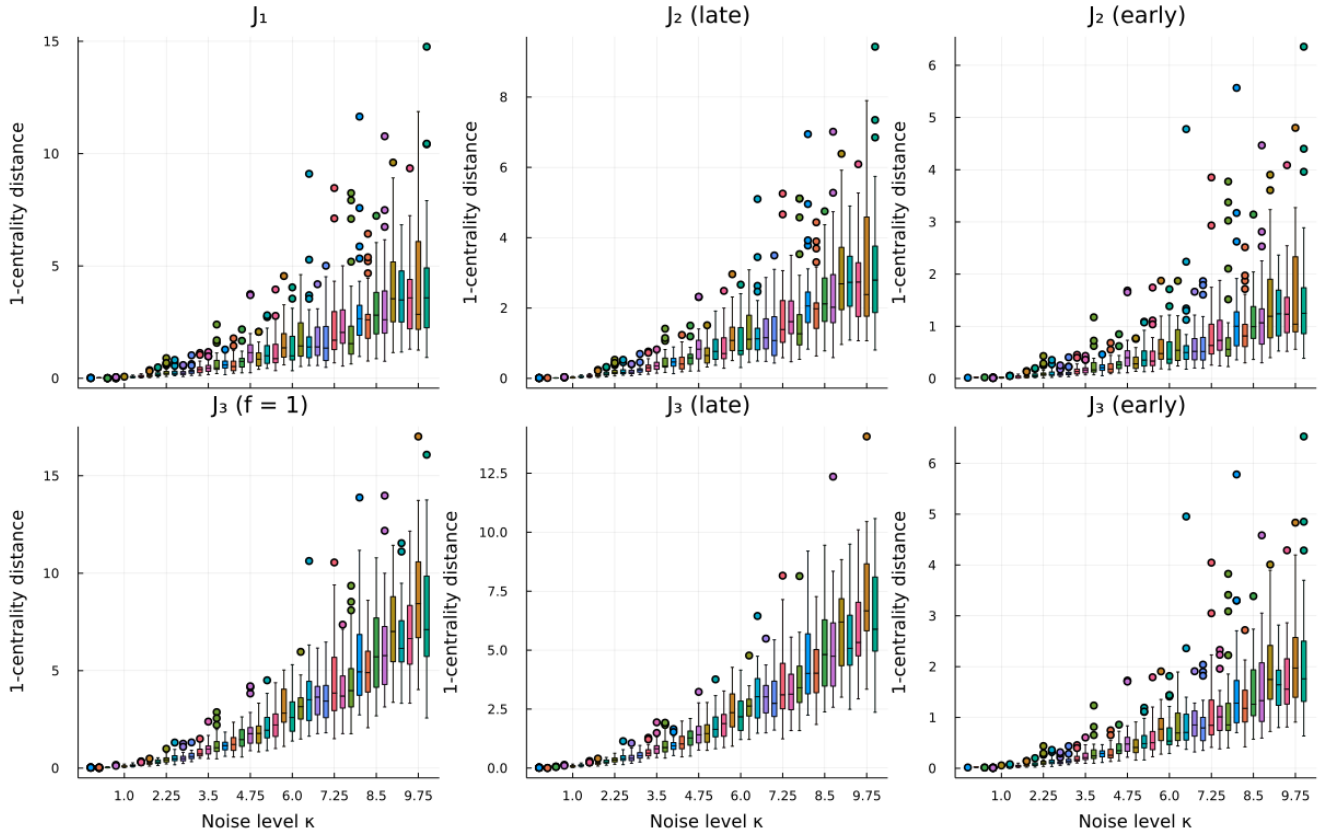


Figure 7: Boxplots for the 1-centrality distance between the centrality functions of the point cloud in Figure 6a and its perturbations for all noise levels.

We now establish bounds for the  $p$ -centrality distance to quantify the stability of the centrality measures. We simplify notations by dropping  $n$  and write  $J_\sigma$  when we consider any of the centrality functions  $J_{1,\sigma}$ ,  $J_{2,\sigma}$ , and  $J_{3,\sigma}$ .

The following results focus on collections  $\Lambda$  and  $\Lambda'$  of persistent homology class representatives obtained through a filtration of a simplicial complex. We define

$$K = \{\max(d(\sigma) - b(\sigma)) : \sigma \in \Lambda \cup \Lambda'\}$$

as the maximum persistence observed across all cycles in both  $\Lambda$  and  $\Lambda'$ . Additionally, let  $q = \max\{|\Lambda|, |\Lambda'|\}$  be the maximum number of cycles between  $\Lambda$  and  $\Lambda'$ . Lastly, we assume that there exists a homology class with a non-empty first-order merge cluster.

**Lemma 3.5.** *Let  $\mathcal{J}(\Lambda)$  be the collection of centrality measures generated from  $\Lambda$ . Then*

$$C_p(\mathcal{J}(\Lambda), \mathcal{J}(\Lambda')) \leq \begin{cases} K^{1+\frac{1}{p}}(1+q), & \text{if } 1 \leq p < \infty, \\ Kq(1+q), & \text{if } p = \infty. \end{cases}$$

*Proof.* Monotonicity of the centrality measures (Proposition 3.1) yields

$$J_\sigma(\epsilon) \leq P_\epsilon(\sigma) + \sum_{\tau} \sum_{[v] \in M_\tau[\sigma, d(\sigma)]} f_\sigma(\phi) P_\epsilon(v).$$

By the definitions of  $K$  and  $q$ , we have  $J_\sigma(\epsilon) \leq K + f_\sigma(\phi)qK$ . Thus, by the monotonicity of Lebesgue integrals, we obtain  $\|J_\sigma(\epsilon)\|_p \leq \sqrt[p]{d^* [K + f_\sigma(\phi)qK]}$ , where the integral is taken from  $[0, d^*]$ . Recall that  $f_\sigma(\phi) \in (0, 1]$ . Thus,

$$\|J_\sigma(\epsilon)\|_p \leq K^{1+\frac{1}{p}}(1+q). \quad (3.2)$$

Now, we define  $X = \{\|J_\sigma\| : J_\sigma \in \mathcal{J}(\Lambda)\}$  and  $Y = \{\|J'_\delta\| : J'_\delta \in \mathcal{J}(\Lambda')\}$ . Assuming  $|\Lambda| \leq |\Lambda'|$ , we consider a bijection  $\phi : X \cup \Delta \rightarrow Y$ . Note that  $|\|J_\sigma\|_p - \|J_\delta\|_p| \leq \max\{\|J_\sigma\|_p, \|J_\delta\|_p\}$ . Using equation (3.2), for any  $x \in X$  and  $y \in Y$ ,  $|x - y| \leq K^{1+\frac{1}{p}}(1+q)$ . Now, suppose  $x \in \Delta$ . By Equation (3.1), we obtain

$$\frac{1}{2} \max\{x, \phi(x)\} \leq \frac{1}{2} K^{1+\frac{1}{p}}(1+q) < K^{1+\frac{1}{p}}(1+q).$$

Therefore,  $C_p(\mathcal{J}(\Lambda), \mathcal{J}(\Lambda')) \leq K^{1+\frac{1}{p}}(1+q)$  for  $1 \leq p < \infty$ .

We consider the case where  $p = \infty$ . Note that, for any  $\epsilon \geq 0$ ,  $J_\sigma(\epsilon) \geq 0$ . Hence,

$$\sum_m \|J_m - J'_m\|_\infty \leq \sum_m \max\{J_m(d^*), J'_m(d^*)\}.$$

Since  $J_\sigma(\epsilon) \leq K + f_\sigma(\phi)qK$  and  $f_\sigma(\phi) \in (0, 1]$ , we have

$$\sum_m \|J_m - J'_m\|_\infty \leq \sum_m K(1+q).$$

By applying the definition of  $q$ , we have proven the lemma.  $\square$

The value of  $K$  suggests a potential for variations in the lifetime of topological features. In the previous lemma, a smaller value of  $K$  contributes to tighter bounds.

We now introduce a constant that will be instrumental in establishing an upper bound for the centrality distance. This constant is given by

$$R(p) := \begin{cases} \sqrt[p]{2}K(1+q), & \text{if } 1 \leq p < \infty, \\ 2q(1+q), & \text{if } p = \infty. \end{cases}$$

The next step is to demonstrate that the  $p$ -centrality distance between any two collections of centrality measures is always upper bounded by this constant and the bottleneck distance between the corresponding collections.

**Theorem 3.6.** *Let  $D$  and  $D'$  represent the persistence diagrams corresponding to the collections  $\Lambda$  and  $\Lambda'$ . Then*

$$C_p(\mathcal{J}(\Lambda), \mathcal{J}(\Lambda')) \leq \begin{cases} R(p) \sqrt[p]{d_B(D, D')}, & \text{if } 1 \leq p < \infty, \\ R(p) d_B(D, D'), & \text{if } p = \infty. \end{cases}$$

*Proof.* For any pair of homology classes  $[\sigma]$  and  $[\delta]$ , application of the triangle inequality yields

$$|P_\epsilon(\sigma) - P_\epsilon(\delta)| \leq |d(\sigma) - d(\delta)| + |b(\sigma) - b(\delta)|.$$

By the definition of the bottleneck distance, we have  $|P_\epsilon(\sigma) - P_\epsilon(\delta)| \leq 2d_B(D, D')$ . For the left-hand expression, the maximum taken over all homology classes  $[\sigma]$  and  $[\delta]$  is  $P_\epsilon(\sigma)$  or  $P_\epsilon(\delta)$ . Consequently,  $K \leq 2d_B(D, D')$ . The conclusion follows from Lemma 3.5.  $\square$

The following corollaries leverage the combinatorial stability theorem [6, p. 123] to reformulate the previously established bounds in terms of the constant  $R(p)$  and properties of the filtration functions. We omit the proofs for these corollaries as they directly apply the referenced theorem.

**Corollary 3.7.** *Let  $w, w' : \mathcal{C} \rightarrow \mathbb{R}$  be monotone real-valued functions that filter the simplicial complex  $\mathcal{C}$ . Then*

$$C_p(\mathcal{J}(\Lambda), \mathcal{J}(\Lambda')) \leq \begin{cases} R(p) \sqrt[p]{\|w - w'\|_\infty}, & \text{if } 1 \leq p < \infty, \\ R(p) \|w - w'\|_\infty, & \text{if } p = \infty. \end{cases}$$



In the next corollary, we introduce another constant given by

$$R'(p) := \begin{cases} \sqrt[p]{2K}(1+q'), & \text{if } 1 \leq p < \infty, \\ 2q'(1+q'), & \text{if } p = \infty, \end{cases}$$

where  $q' := \max\{\sum_r |M_r[\sigma, d(\sigma)]| : \sigma \in \Lambda \cup \Lambda' \text{ and } P_e(\sigma) \neq 0\}$ . The term  $q'$  represents the maximum number of successive merging of homology classes across  $\Lambda$  and  $\Lambda'$ .

**Corollary 3.8.** *Let  $D$  and  $D'$  be persistence diagrams corresponding to the collections  $\Lambda$  and  $\Lambda'$ . Then*

$$C_p(\mathcal{J}_n(\Lambda), \mathcal{J}_n(\Lambda')) \leq \begin{cases} R'(p) \sqrt[p]{d_B(D, D')}, & \text{if } 1 \leq p < \infty, \\ R'(p) d_B(D, D'), & \text{if } p = \infty. \end{cases}$$

The previous corollary refines the bounds further by incorporating information about the merge clusters of homology classes. A smaller  $q'$  indicates less variation in how homology classes merge between  $\Lambda$  and  $\Lambda'$ , leading to potentially tighter bounds.

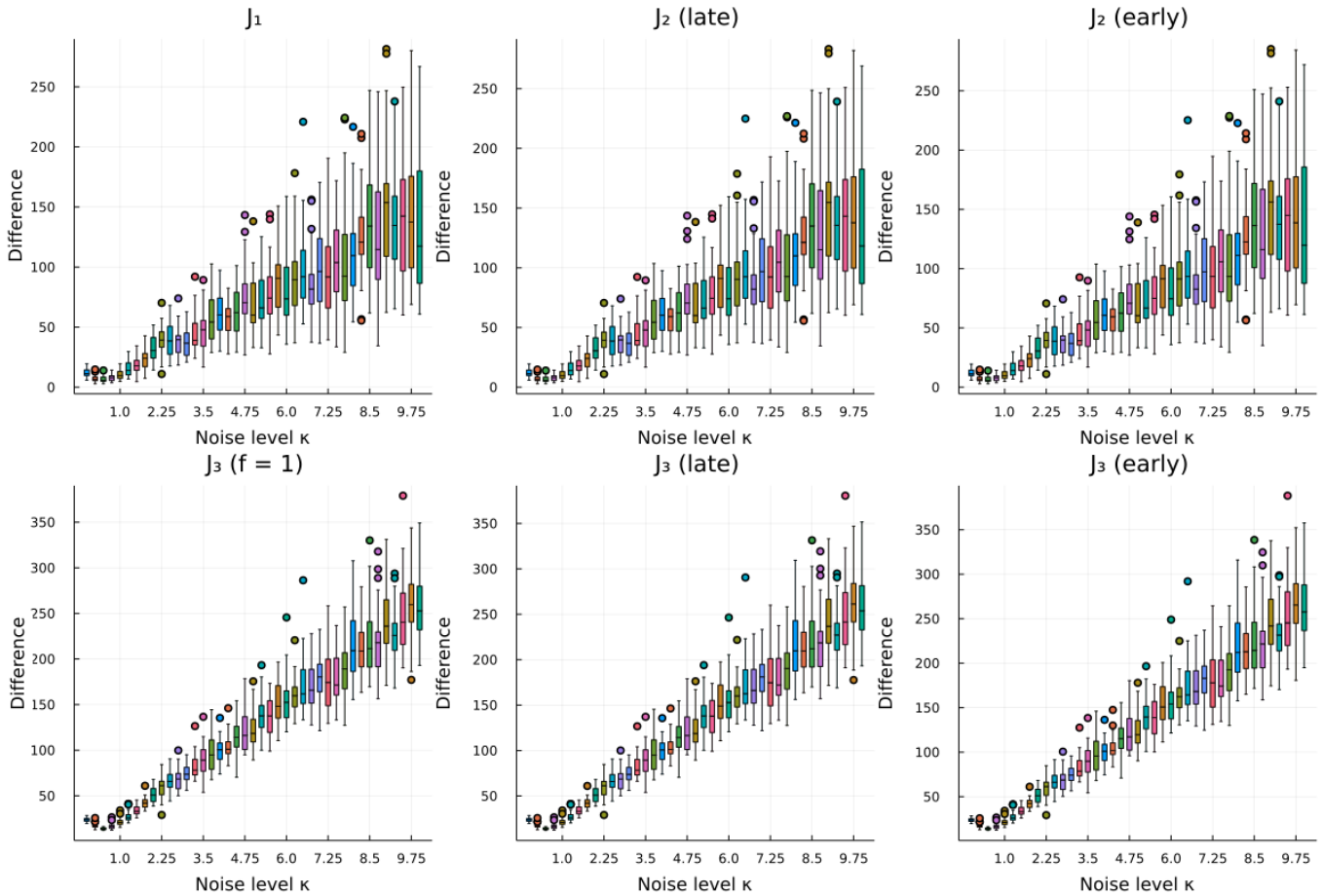


Figure 8: Boxplots of the difference between the 1-centrality distance in Figure 7 and corresponding bounds given by Proposition 3.8.

**Theorem 3.9 (Stability).** *Let  $w, w' : \mathcal{C} \rightarrow \mathbb{R}$  be monotone real-valued functions that filter the simplicial complex  $\mathcal{C}$ . The inequality*

$$C_p(\mathcal{J}_n(\Lambda), \mathcal{J}_n(\Lambda')) \leq R'(p) \|w - w'\|_\infty$$

*holds when either  $p = \infty$ , or  $1 \leq p < \infty$  and  $q' > (1/\sqrt[p]{2K}) - 1$ .*

*Proof.* Suppose that  $1 \leq p < \infty$  and  $q' > \left(1/\sqrt[p]{2K}\right) - 1$ . It follows from the combinatorial stability theorem and Corollary 3.8 that

$$\frac{1}{R'(p)} C_p(\mathcal{J}_n(\Lambda), \mathcal{J}_n(\Lambda')) \leq \|w - w'\|_{\infty}^{\frac{1}{p}}.$$

The conclusion follows because  $x^{\frac{1}{p}} \leq x$  for any  $x \geq 0$ .

Suppose  $p = \infty$ . The inequality  $2q'(1 + q') \geq 1$  holds since  $q' \geq 1$ . Applying the combinatorial stability theorem, we obtain

$$\frac{1}{R'(p)} C_p(\mathcal{J}_n(\Lambda), \mathcal{J}_n(\Lambda')) \leq \|w - w'\|_{\infty}.$$

□

To evaluate the effectiveness of the bounds established in Corollary 3.8, we perform the following analysis. We consider the 1-centrality distances (refer to Example 3.4) calculated for various perturbation levels. From the bounds provided by Corollary 3.8, we subtract the actual 1-centrality distances. These differences are then plotted as boxplots across increasing levels of perturbations, similar to the approach used in the previous figures. This visualization allows us to assess how closely the theoretical bounds align with the empirical observations of the 1-centrality distance under increasing network perturbations.

#### 4. Application to fractal-like point clouds

This chapter explores the application of centrality measures to self-similar point clouds. While toy datasets offer valuable starting points, their applicability to real-world scenarios may be limited. In contrast, fractal-like point clouds, with their inherent complexity and potential for higher dimensionality, have been vastly documented to appear in nature [17]. Persistent homology has been studied on fractals such as computing the affine fractals from landscapes [5] and estimating the fractal dimension [16]. Analyzing these structures allows us to explore the generalizability and effectiveness of centrality measures in deciphering intricate network-like relationships within spatial data.

For this application, we employ a method for separating signal from noise in persistence diagrams using a hypothesis testing approach developed by Bobrowski and Skraba [1]. This method relies on the concept of *multiplicative persistence* ( $\pi(p) = \frac{d}{b}$ ) for any birth-death pair  $(b, d)$  in the  $k^{\text{th}}$  persistence diagram  $\text{dgm}_k$ . Throughout this chapter, we operate under the assumption of the validity of the strong universality conjecture. Within this theoretical framework, a left-skewed Gumbel distribution (LGumbel) plays a pivotal role in the hypothesis testing process. In the conjecture below, the constant  $\lambda$  is the Euler-Mascheroni constant. Moreover,  $A(\mathcal{T})$  is 1 if  $\mathcal{T}$  is the Rips filtration, and 0.5 if  $\mathcal{T}$  is the Čech filtration.

**Conjecture 4.1.** Consider  $d$ -dimensional metric measure space  $S$  and a sequence of random variables  $\mathbf{X}_n = (X_1, \dots, X_n) \in S^n$  with joint probability law  $\mathbb{P}_n$ . Let  $\mathcal{S} = (S, \mathbb{P})$  be a sampling model under a filtration type  $\mathcal{T}$ . For any  $\mathcal{S} \in \mathcal{U}, \mathcal{T}$ , and  $k \geq 1$ , the limit of  $\mathcal{L}_n$  as  $n$  approaches infinity equals the left-skewed Gumbel distribution, where

$$\mathcal{L}_n(\mathcal{S}, \mathcal{T}, k) = \frac{1}{|\text{dgm}_k|} \sum_{p \in \text{dgm}_k} \delta_{l(p)}$$

and  $l(p) = A(\mathcal{T}) \log \log(\pi(p)) - \lambda - \bar{L}$ . The expression  $\bar{L}$  is given by

$$\frac{1}{|\text{dgm}_k|} \sum_{p \in \text{dgm}_k} \log \log(\pi(p)).$$

This conjecture states that, as the sample size approaches infinity, the distribution of certain features derived from persistence diagrams converges to the left skewed Gumbel (LGumbel) distribution.

From the hypothesis testing framework, we recover the signal part  $\text{dgm}_k^S$  of the diagram by considering only those  $p \in \text{dgm}_k$ , where

$$e^{-e^{l(p)}} < \frac{\alpha}{|\text{dgm}_k|}$$

for a significance level  $\alpha$ . Until the end of this chapter, we consider the significance level  $\alpha = 0.05$ .

We now explore the topological structure of a point cloud containing 1,000 points sampled around the well-known Sierpinski Triangle (Figure 9). Given the inherent variability in point cloud data and the complex nature of TDA, we employ a bootstrapping approach to assess the robustness of our findings regarding the number of holes (signals) identified.

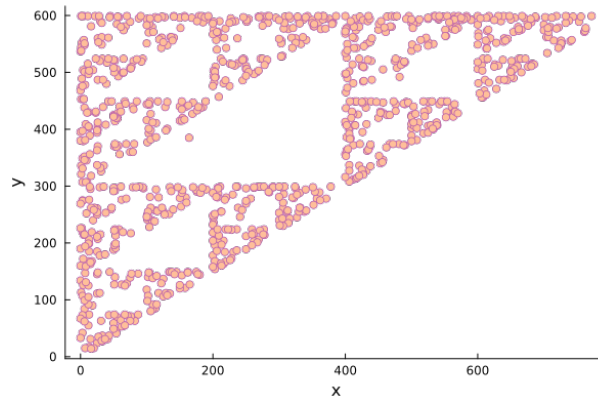


Figure 9: A point cloud sampled around a Sierpinski triangle.

Bootstrapping [7] involves creating multiple random samples (with replacement) from the original point cloud. In this case, we generated 1,000 samples, each containing 800 points (80% of the original data). This sampling percentage is chosen following the rule of thumb that using at least 50% of the data is beneficial for obtaining statistically significant results in bootstrapping.

We then use the hypothesis testing framework for each bootstrapped sample to identify the number of signals. The mean number of signals across all bootstrapped samples is 0.68, with an approximate standard error of 0.018. The 95% confidence interval for the number of holes is (0.643, 0.717). The confidence interval suggests that the true number of signals in the original point cloud likely falls between 0 and 1. However, the interval is skewed slightly closer to 1, indicating a higher probability of there being a single signal present in the data.

We examine one bootstrapped sample whose plot is shown in Figure 10a. Applying the hypothesis testing framework, we identify only a single hole as statistically significant (signal). The corresponding signal is visualized in Figure 10b. The signal contains points from the largest triangle in the fractal.

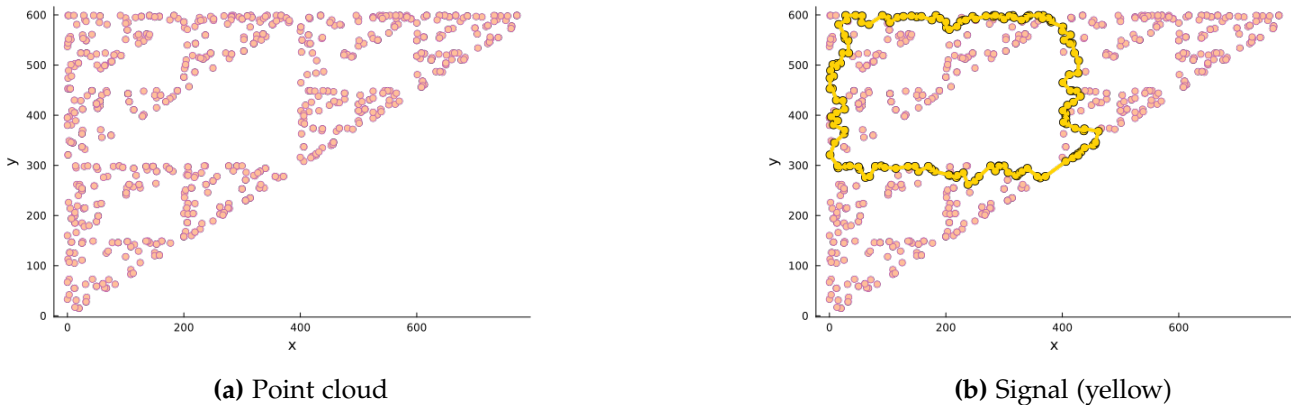


Figure 10: Signal from a bootstrap sample of the point cloud in Figure 9.

In the context of centrality, we consider two primary ideas for analysis. Direct merging events between holes offer insights into the localized dynamics around the merged region. However, a more comprehensive perspective is gained by considering the cumulative effect of merging across the entire filtration process. This broader analysis, encompassing all merging events, can reveal structures with global importance in shaping the overall landscape of significant holes. Next, while centrality measures rank the holes, a quantitative method to pinpoint the most important ones might not be readily available. Here, we address this challenge by examining the relative difference in centrality values between the top-ranked holes and the majority of others. Significant differences in centrality would suggest a higher likelihood of those top-ranked holes being truly central features within the data.

Figure 11 depicts the centrality plots associated with the point cloud. To distinguish between our different centrality measures, we denote centrality functions  $J_3$  with various scaling factors as  $J_4$ ,  $J_5$ , and  $J_6$ . The function  $J_4$  has factors equal to one, while the functions  $J_4$  and  $J_5$  prioritize late merging and early merging, respectively.

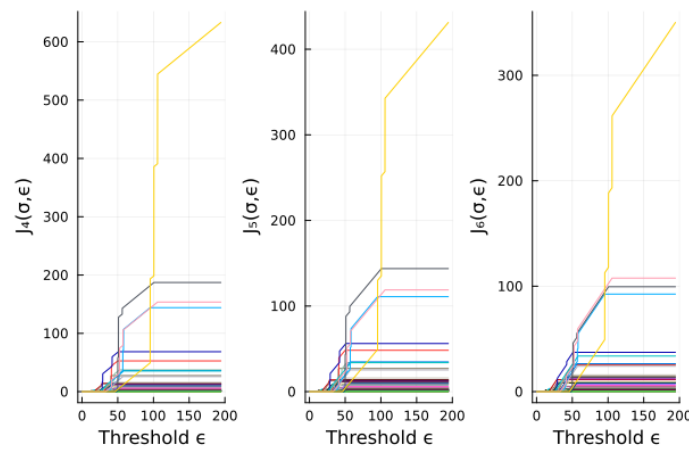


Figure 11: Centrality plots of dimension 1 associated to the Rips filtration of the point cloud in Figure 9.

The centrality plots suggest the presence of a relatively important signal within the point cloud data. This is evidenced by the large difference in the maximum centrality value for the highest ranked hole compared to the others. Notably, this hole coincides with the one previously identified using the hypothesis testing framework. This strong agreement between centrality measures and the well-established topological tools found in persistence diagrams reinforces the significance of this particular hole as a key topological feature within the data.

Further bolstering this observation, we can consider persistence values as a ranking system for holes. Computing the Spearman rank correlation coefficient [18] between the maximum centrality values of  $J_5$  and the persistence values yields a value of 0.997. This value indicates a near-perfect, monotonically increasing relationship between the two rankings. In simpler terms, holes ranked highly by centrality also tend to have high persistence values. By aligning with established methods like persistence diagrams, these findings suggest that centrality functions effectively capture features similar to those identified by common TDA summaries.

We examine other bootstrap sample (Figure 12a), where the hypothesis testing framework fails to identify a signal. Interestingly, the persistence diagram (Figure 12c) exhibits a point far from the diagonal, potentially indicating a feature. However, this point is not classified as a signal by the hypothesis testing framework.

Despite the lack of a signal using traditional methods, both the persistence and the centrality functions, viewed as rankings, reveal some patterns. Mirroring the previous bootstrap sample, the highest ranked hole by both methods corresponds to points from the largest triangle in the fractal. Similarly, the second highest ranked hole aligns with points from the next largest triangle.

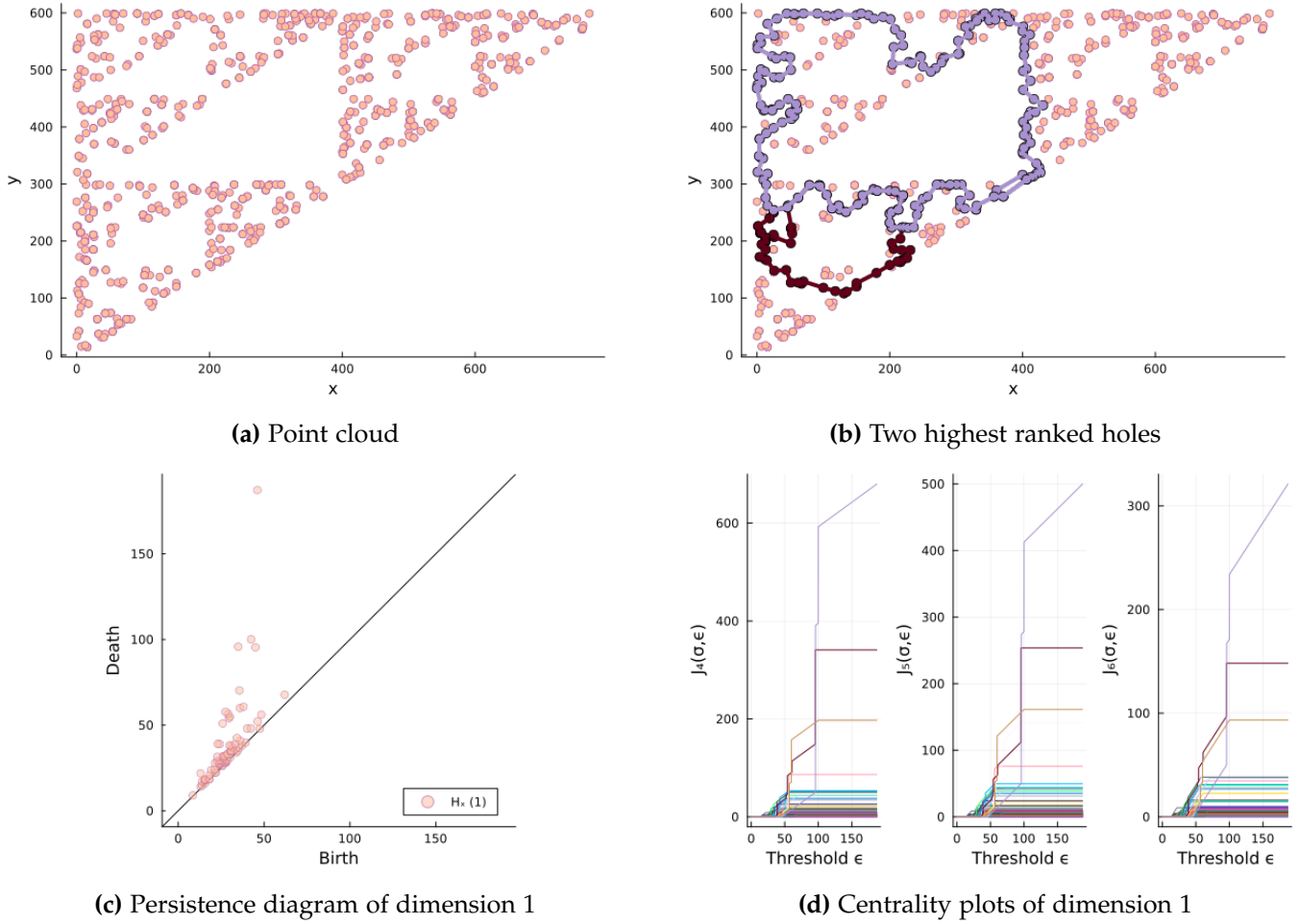
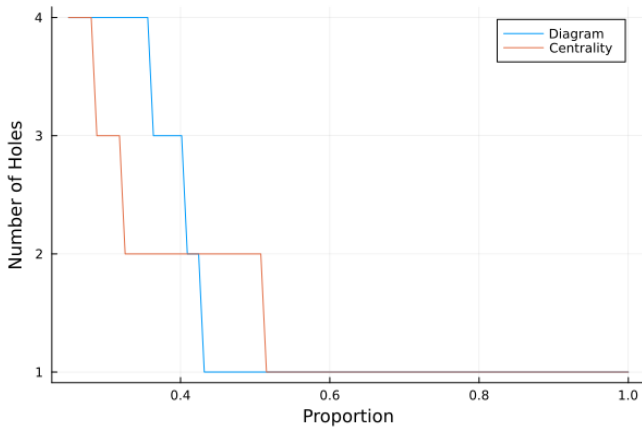


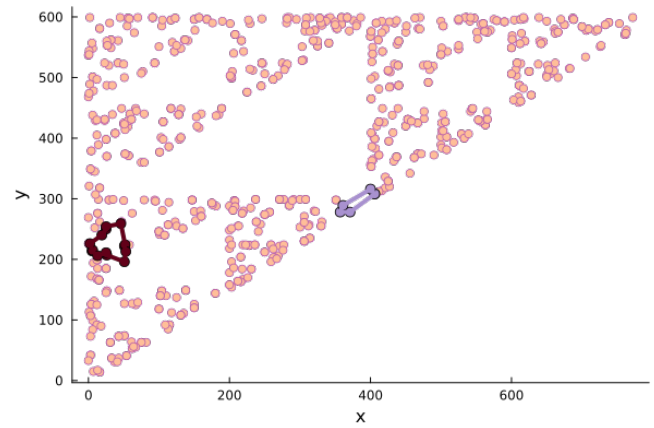
Figure 12: Two highest ranked holes according to the persistence diagram and the centrality plots from a bootstrap sample of the point cloud in Figure 9 with no identified signals.

To gain a deeper understanding, we delve into the values of both summaries and how they identify holes. We examine how many holes satisfy a threshold, defined as a percentage (i) of the maximum value from each summary. We will analyze a range of percentages for i from 0.25 to 1. This analysis will be performed for both the persistence diagram and the centrality function  $J_5$ . Figure 13a illustrates the graph of the number of holes identified by each method at different thresholds. The graph reveals interesting differences in how persistence and centrality identify holes. While  $J_5$  identifies two holes meeting the threshold at  $i = 0.5$ , the persistence diagram identifies a single hole at a slightly lower threshold (around 0.4). As we move to lower thresholds, the persistence diagram consistently identifies more holes than  $J_5$  at each step. These observations suggest that centrality has the tendency to be selective in identifying features compared to the persistence diagram. These observations highlight the potential of centrality functions to capture features that might differ when using common topological summaries.

We can also identify and analyze the merging events across the entire filtration process. Figure 13b shows the earliest homology class considering all possible orders of merging clusters in each of the two highest ranked holes. We note the small diameter of each hole. This observation suggests that some short-lived holes, presented during the early stages of filtration, contribute to the overall centrality of the larger holes. Additionally, the centrality of these large holes appears to be influenced by merge instances originating from them much earlier in the filtration process.

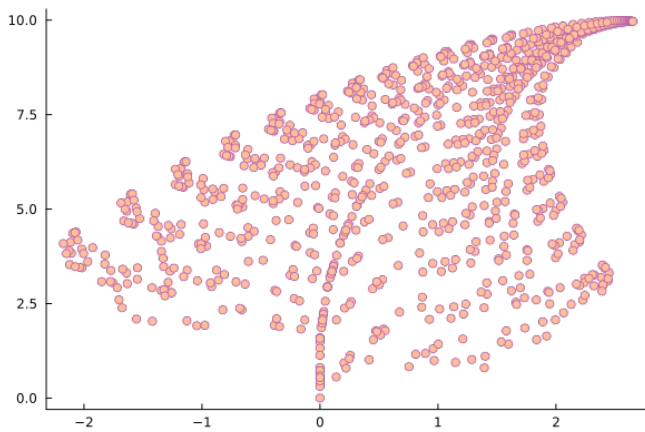


(a) Number of holes satisfying the threshold

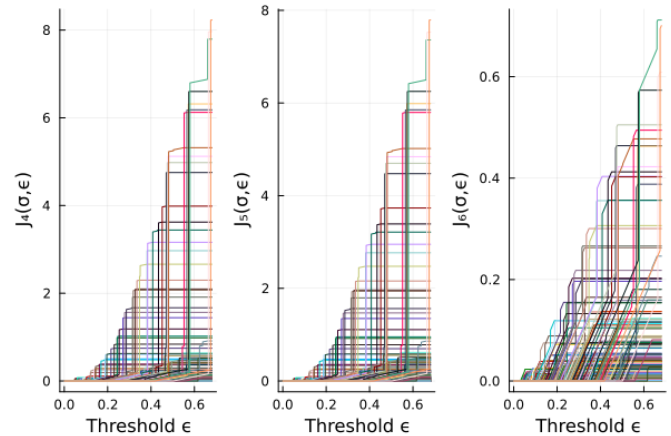


(b) Earliest cycles across all orders of merge clusters

Figure 13: Additional features that the centrality functions capture from the point cloud in Figure 12a.



(a) Point cloud



(b) Centrality plots of dimension 1

Figure 14: Centrality plots of dimension 1, produced by the Rips filtration of a point cloud sampled around the Barnsley Fern.

We conclude this section by examining a different point cloud, sampled around the Barnsley Fern fractal (Figure 14a). Here, we employ the same bootstrapping approach used previously. No signals were identified in any of the bootstrapped samples for the Barnsley Fern. The Spearman rank correlation between the maximum  $J_5$  centrality values and the persistence values remains high (0.893), indicating a strong monotonically increasing relationship. Similar to the previous example, this suggests that centrality and persistence rankings tend to agree on the relative importance of holes despite the absence of identifiable signals. Furthermore, Figure 15a shows the three highest ranked holes identified by centrality. Extracting the earliest cycles across all possible orders of merge clusters for these three holes reveals a single cycle. In simpler terms, this implies that one hole merged with another, then the survivor merged with the third hole, forming a single connected component at some point during the filtration process. The representative of the single cycle further suggests that the corresponding hole might have played a central role in connecting the other two holes during the filtration process.



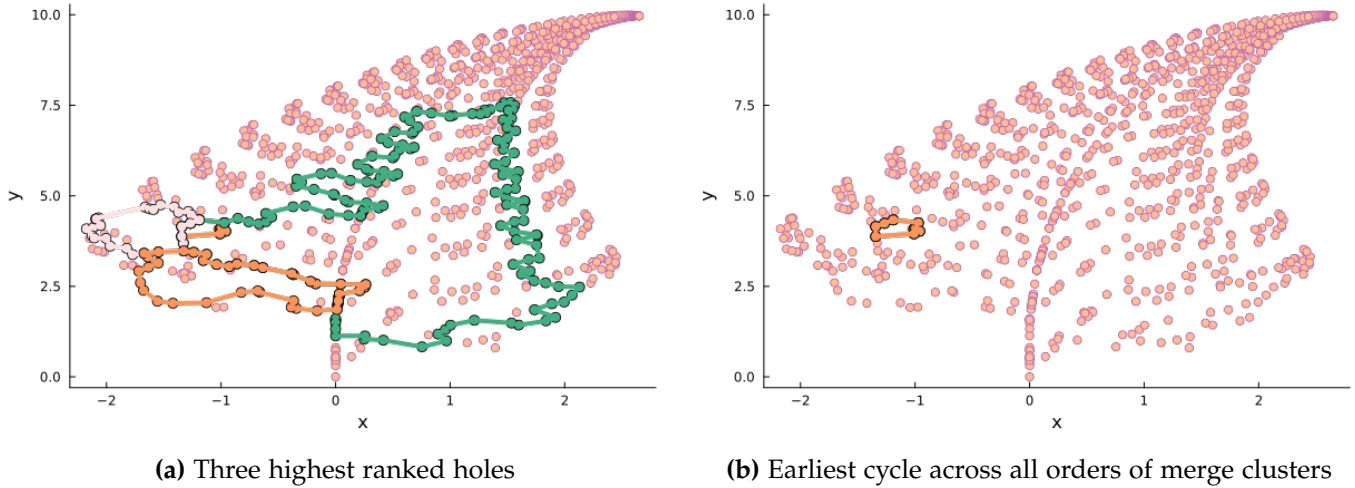


Figure 15: The earliest cycle across all orders of merge clusters of each of the three highest ranked holes

Similar to our previous analysis, we compared how centrality and persistence identify holes based on thresholds. Figure 16 illustrates the number of holes detected by each method at different threshold levels (represented as a percentage of the maximum persistence value). The plot reveals key differences in how these methods classify holes. At higher thresholds ( $i > 0.6$ ), the centrality function  $J_5$  identifies holes that meet the threshold. However, the persistence diagram identifies a larger number of holes, including those with lower persistence values ( $i \leq 0.6$ ). This trend continues as we move towards lower thresholds. The persistence diagram consistently detects a higher number of holes (at least 50 holes at  $i = 0.25$ ) compared to the centrality plot (which identifies at most 20 holes at  $i = 0.25$ ). These observations reinforce our earlier finding that the centrality functions might be more selective in identifying features depending on a certain threshold.

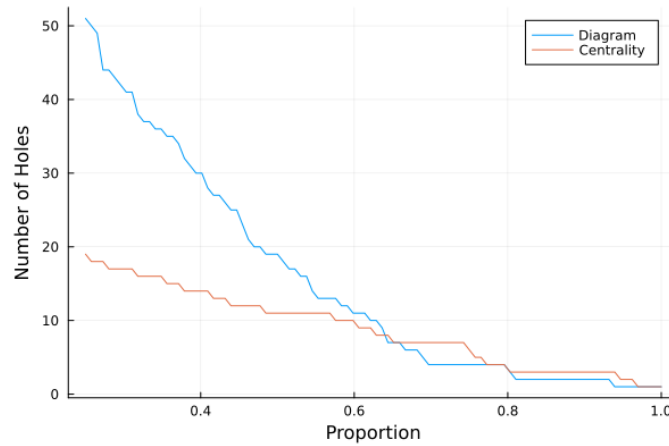


Figure 16: Number of holes satisfying the threshold depending on the proportion of the maximum persistence and the maximum  $J_5$  centrality values.

## 5. Conclusion

We introduced novel centrality measures that leverage both persistence and merge dynamics of homology classes. These measures aim to capture a more comprehensive picture of the topological structure within point cloud data compared to traditional summaries. The algorithm for computing the merge dynamics of homology classes is guided by the equivalence of merging and  $q$ -nearness between two classes. Similar to persistence barcodes, we have generated shape descriptors in the form of plots and heat maps and established stability by defining a pseudo-metric similar to bottleneck and landscape distances. Our

initial investigation with self-similar point clouds has demonstrated agreement with existing TDA tools, while also revealing additional properties of important features.

Moving forward, we plan to assess the efficacy of these measures across diverse real-world point cloud datasets and within machine learning contexts. Furthermore, future research endeavors will focus on refining the centrality functions and investigating their mathematical properties in greater depth.

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