

Computing the edge metric dimension of convex polytopes related graphs



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Abstract

Let $G = (V(G), E(G))$ be a connected graph and $d(f, y)$ denotes the distance between edge f and vertex y , which is defined as $d(f, y) = \min\{d(p, y), d(q, y)\}$, where $f = pq$. A subset $W_E \subseteq V(G)$ is called an edge metric generator for graph G if for every two distinct edges $f_1, f_2 \in E(G)$, there exists a vertex $y \in W_E$ such that $d(f_1, y) \neq d(f_2, y)$. An edge metric generator with minimum number of vertices is called an edge metric basis for graph G and the cardinality of an edge metric basis is called the edge metric dimension represented by $\text{edim}(G)$. In this paper, we study the edge metric dimension of flower graph $f_{n \times 3}$ and also calculate the edge metric dimension of the prism related graphs D'_n and D_n^t . It has been concluded that the edge metric dimension of D'_n is bounded, while of $f_{n \times 3}$ and D_n^t is unbounded.

Keywords: Edge metric dimension, edge metric generator, edge metric basis, resolving set, prism related graphs, flower graph.
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1. Introduction and preliminaries

Slater proposed the idea of metric dimension to find the location of intruder in a network (see [18, 23]). The proposed idea was further extended by Melter and Harary in [10]. Metric dimension is important in robot navigation, chemistry, problems of image processing and pattern recognition etc. (see [4, 5, 11, 12, 14, 15]). The use of metric dimension of graphs was also observed in games like mastermind and coin weighing (see [3]). Cáceres et al. computed the metric dimension of cartesian product of graphs and Hallaway et al. had worked on the metric dimension of permutation of graphs (see [3, 9]). Further Zhang had worked on the theory and resolvability of graphs in [6].

Suppose G is connected graph having edge set $E(G)$ and vertex set $V(G)$, also $|E(G)|$ shows the size of graph G and $|V(G)|$ represents the order of graph G . Let $N(a) = \{b \in V(G) | ab \in E(G)\}$ denote the neighborhood of the vertex a , then $|N(a)|$ is called the degree of the vertex a . Moreover, $\Delta(G)$ and $\delta(G)$ represent the maximum and minimum degree of graph G , respectively.

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The distance between two distinct vertices s and t , is the number of edges in a shortest path between them and it is denoted by $d(s, t)$. A vertex $u \in V(G)$ is said to distinguish a pair of vertices $s, t \in V(G)$ if $d(s, u) \neq d(t, u)$. A set $W \subseteq V(G)$ is a metric generator for G if every pair of vertices of G can be distinguished by some vertex in W . A metric basis is the minimum metric generator for graph G and number of elements in metric basis is called the metric dimension of G , denoted by $\dim(G)$. It was shown that computing the metric dimension of a graph is NP-hard [14].

The edge metric dimension is introduced recently. The concept was brought by Kelenc et al. and further studied by Zubrilina, Peterin, Kratica, Yuezhong Zhang and Ahsan et al. [1, 7, 13, 16, 22, 23]. We can find the distance between an edge $f = pq$ and a vertex y as follows:

$$d(f, y) = \min\{d(p, y), d(q, y)\}.$$

A vertex $a \in V(G)$ is said to distinguish two distinct edges $e_1, e_2 \in E(G)$ if $d(e_1, a) \neq d(e_2, a)$. A set W_E is an edge metric generator of a graph G if every two distinct edges are distinguished by some vertex of W_E . An edge metric basis is the minimum edge metric generator of graph G and its cardinality is called edge metric dimension, denoted by $\text{edim}(G)$.

For an ordered subset $W_E = \{a_1, a_2, \dots, a_k\}$ of the vertex set of $V(G)$, the k -tuple $r(e|W_E) = (d(e, a_1), d(e, a_2), \dots, d(e, a_k))$ is called the edge metric representation of an edge e with respect to W_E . In this sense, W_E is an edge metric generator for G if and only if for every pair of different edges e_1, e_2 of $E(G)$, we have $r(e_1|W_E) \neq r(e_2|W_E)$.

In this whole paper, all vertex indices are considered to be module n . The propositions given below are very helpful for calculating the edge metric dimension of graphs.

Proposition 1.1 ([13]). *If G is a connected graph, then $\text{edim}(G) \geq \lceil \log_2 \Delta(G) \rceil$.*

Proposition 1.2 ([13]). *If G is a connected graph, then $\text{edim}(G) \geq 1 + \lceil \log_2 \delta(G) \rceil$.*

In following proposition, we will demonstrate the metric dimension of flower graph.

Proposition 1.3 ([11]). *For the flower graph $f_{n \times 3}$ with $n \geq 6$, we have*

$$\dim(f_{n \times 3}) = \begin{cases} 2, & n \text{ is even,} \\ 3, & \text{otherwise.} \end{cases}$$

In the following propositions we calculate the metric dimensions of D'_n and D_n^t by showing its resolving sets and the results are obvious.

Proposition 1.4. *For the prism related graph D'_n with $n \geq 4$, we have $\dim(D'_n) = 3$ and its metric basis is $W = \{a_1, a_2, a_{k+2}\}$, where either $n = 2k$ or $n = 2k + 1$.*

Proposition 1.5. *For the prism related graph D_n^t with $n \geq 4$, we have $\dim(D_n^t) = 3$ and its metric basis is $W = \{a_1, a_2, a_{k+1}\}$, where either $n = 2k$ or $n = 2k + 1$.*

The rest of paper is structured as follows. In Section 2, edge metric dimension of flower graph $f_{n \times 3}$ will be studied. In Section 3, edge metric dimension of prism related graph D'_n will be investigated. In Section 4, edge metric dimension of prism related graph D_n^t will be determined. In last Section, article will be concluded.

2. Edge metric dimension of flower graph

In this Section, we will investigate the edge metric dimension of flower graph $f_{n \times 3}$. We have the vertex set $V(f_{n \times 3}) = \{a_\alpha, b_\alpha | 1 \leq \alpha \leq n\}$ and the edge set $E(f_{n \times 3}) = \{a_\alpha a_{\alpha+1}, a_\alpha b_\alpha, a_\alpha b_{\alpha+1} | 1 \leq \alpha \leq n\}$ as shown in Figure 1.

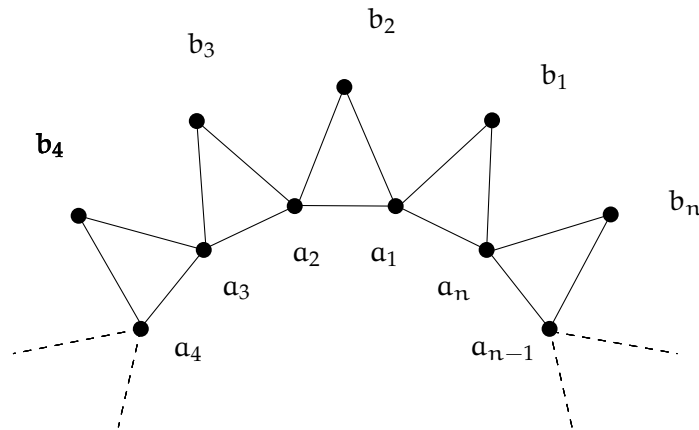


Figure 1: Graph of $f_{n \times 3}$.

Lemma 2.1. Let $Y = \{b_1, b_2, \dots, b_n\}$ be a subset of $V(f_{n \times 3})$. Then any arbitrary edge metric generator W_E of $f_{n \times 3}$ contains at least $\lceil \frac{n}{2} \rceil$ vertices of Y .

Proof. Suppose that W_E contains at most $\lceil \frac{n}{2} \rceil - 1$ vertices of Y for a contradiction. Without loss of generality we assume that $b_\alpha, b_{\alpha+1} \notin W_E$, then we have $(a_\alpha b_\alpha | W) = (a_\alpha b_{\alpha+1} | W)$, a contradiction. \square

Remark 2.2. Let W_E be any edge metric basis of $f_{n \times 3}$. We note that W_E contains all odd vertices (vertex indices are odd) of Y for odd n , while W_E contains either all odd vertices or even vertices (vertex indices are even) of Y for even n .

Theorem 2.3. For the flower graph $f_{n \times 3}$, we have

$$\text{edim}(f_{n \times 3}) = \begin{cases} 3, & n = 4, \\ 4, & n = 3, 5, \\ \lceil \frac{n}{2} \rceil, & \text{otherwise.} \end{cases}$$

Proof. For $n = 3, 4, 5$ we have calculated the edge metric dimension by total enumeration, $\text{edim}(f_{3 \times 3}) = 4$ and its edge metric basis is $\{a_1, a_2, b_1, b_2\}$, $\text{edim}(f_{4 \times 3}) = 3$ and its edge metric basis is $\{b_1, b_3, b_4\}$, and $\text{edim}(f_{5 \times 3}) = 4$ and its edge metric basis is $\{a_1, b_1, b_3, b_5\}$.

For $n \geq 6$, we discuss the following four cases.

Let $W_E = \{b_1, b_3, b_5, \dots, b_{n-1}\}$. We will show that W_E is an edge metric basis of $f_{n \times 3}$ in Case(I) and Case(II), respectively.

Case (I) When $n \equiv 0 \pmod{4}$. We can write $n = 2k$, $k \geq 4$, and k is even. Let $W_1 = \{b_1, b_3, b_{k+1}, b_{k+3}\}$. Next, we give representation of edges of $f_{n \times 3}$ with respect to W_1 .

$$r(a_\alpha a_{\alpha+1} | W_1) = \begin{cases} (\alpha, 1, k - \alpha, \alpha + k - 2), & 1 \leq \alpha \leq 2, \\ (\alpha, \alpha - 2, k - \alpha, k - \alpha + 2), & 3 \leq \alpha \leq k - 1, \\ (k, k - 2, 1, 2), & \alpha = k, \\ (n - \alpha, \alpha - 2, \alpha - k, 1), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha, n - \alpha + 2, \alpha - k, \alpha - k - 2), & k + 3 \leq \alpha \leq n - 1, \\ (1, 2, k, k - 2), & \alpha = n, \end{cases}$$

$$r(a_\alpha b_\alpha | W_1) = \begin{cases} (0, 2, k, k-1), & \alpha = 1, \\ (\alpha, 3-\alpha, k-\alpha+1, k), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha-2, k-\alpha+1, k-\alpha+3), & 4 \leq \alpha \leq k, \\ (n-k, k-1, 0, 2), & \alpha = k+1, \\ (n-k-1, k, 2, 1), & \alpha = k+2, \\ (n-k-2, n-k, 3, 0), & \alpha = k+3, \\ (n-\alpha+1, n-\alpha+3, \alpha-k, \alpha-k-2), & k+4 \leq \alpha \leq n, \end{cases}$$

$$r(a_\alpha b_{\alpha+1} | W_1) = \begin{cases} (1, 2, k, k-1), & \alpha = 1, \\ (2, 0, k-1, k), & \alpha = 2, \\ (\alpha, \alpha-2, k-\alpha+1, k-\alpha+3), & 3 \leq \alpha \leq k-1, \\ (k, k-2, 0, 3), & \alpha = k, \\ (n-k, k-1, 1, 2), & \alpha = k+1, \\ (n-k-1, k, 2, 0), & \alpha = k+2, \\ (n-\alpha+1, n-\alpha+3, \alpha-k, \alpha-k-2), & k+3 \leq \alpha \leq n-1, \\ (0, 3, k, k-2), & \alpha = n. \end{cases}$$

From above representation we see that $r(a_\alpha b_\alpha | W_1) = r(a_\alpha b_{\alpha+1} | W_1)$ when $1 \leq \alpha \leq n$ and $\alpha \neq 1, 2, 3, k, k+1, k+2, k+3, n$ and no other edges have same representation. If we take odd α , where $1 \leq \alpha \leq n$ and $\alpha \neq 1, 3, k+1, n-1$ such that $W_E = W_1 \cup \{b_\alpha\}$ then $r(a_\alpha b_\alpha | W_E) \neq r(a_\alpha b_{\alpha+1} | W_E)$ which implies that $\text{edim}(f_{n \times 3}) \leq \lceil \frac{n}{2} \rceil$. So from Lemma 2.1, W_E is an edge metric generator for $f_{n \times 3}$ and $\text{edim}(f_{n \times 3}) = \lceil \frac{n}{2} \rceil$.

Case (II) When $n \equiv 2 \pmod{4}$. We can write $n = 2k, k \geq 3$, and k is odd. Let $W_1 = \{b_1, b_3, b_{k+2}\}$. Next, we give representation of edges of $f_{n \times 3}$ with respect to W_1 .

$$r(a_\alpha a_{\alpha+1} | W_1) = \begin{cases} (\alpha, 1, k-\alpha+1), & 1 \leq \alpha \leq 2, \\ (\alpha, \alpha-2, k-\alpha+1), & 3 \leq \alpha \leq k, \\ (n-k-1, k-1, 1), & \alpha = k+1, \\ (n-\alpha, n-\alpha+2, \alpha-k-1), & k+2 \leq \alpha \leq n-1, \\ (1, 2, n-k-1), & \alpha = n, \end{cases}$$

$$r(a_\alpha b_\alpha | W_1) = \begin{cases} (0, 2, k), & \alpha = 1, \\ (\alpha, 3-\alpha, k-\alpha+2), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha-2, k-\alpha+2), & 4 \leq \alpha \leq k, \\ (n-\alpha+1, \alpha-2, k-\alpha+2), & k+1 \leq \alpha \leq k+2, \\ (n-\alpha+1, n-\alpha+3, \alpha-k-1), & k+3 \leq \alpha \leq n, \end{cases}$$

$$r(a_\alpha b_{\alpha+1} | W_1) = \begin{cases} (1, 2, k), & \alpha = 1, \\ (\alpha, \alpha-2, k-\alpha+2), & 2 \leq \alpha \leq k, \\ (n-\alpha+1, \alpha-2, \alpha-k-1), & k+1 \leq \alpha \leq k+2, \\ (n-\alpha+1, n-\alpha+3, \alpha-k-1), & k+3 \leq \alpha \leq n-1, \\ (0, 3, k-1), & \alpha = n. \end{cases}$$

From above representation we see that $r(a_\alpha b_\alpha | W_1) = r(a_\alpha b_{\alpha+1} | W_1)$ when $1 \leq \alpha \leq n$ and $\alpha \neq 1, 2, 3, k+1, k+2, n$ and no other edges have same representation. If we take odd α , where $1 \leq \alpha \leq n$ and $\alpha \neq 1, 3, k+2$ such that $W_E = W_1 \cup \{b_\alpha\}$ then $r(a_\alpha b_\alpha | W_E) \neq r(a_\alpha b_{\alpha+1} | W_E)$ which implies that $\text{edim}(f_{n \times 3}) \leq \lceil \frac{n}{2} \rceil$. So from Lemma 2.1, W_E is an edge metric generator for $f_{n \times 3}$ and $\text{edim}(f_{n \times 3}) = \lceil \frac{n}{2} \rceil$.

Let $W_E = \{b_1, b_3, b_5, \dots, b_{n-2}, b_n\}$. We will show that W_E is an edge metric basis of $f_{n \times 3}$ in Cases (III) and (IV), respectively.

Case (III) When $n \equiv 1 \pmod{4}$. We can write $n = 2k + 1$, $k \geq 4$, and k is even. Let $W_1 = \{b_1, b_3, b_{k+1}, b_{k+3}, b_{k+5}\}$. Next, we give representation of edges of $f_{n \times 3}$ with respect to W_1 .

$$\begin{aligned}
 r(a_\alpha a_{\alpha+1} | W_1) &= \begin{cases} (1, 1, k-1, k, k-2), & \alpha = 1, \\ (\alpha, 1, k-\alpha, k-\alpha+2, \alpha+k-3), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha-2, k-\alpha, k-\alpha+2, k-\alpha+4), & 4 \leq \alpha \leq k-1, \\ (k, k-2, 1, 2, 4), & \alpha = k, \\ (n-\alpha, \alpha-2, \alpha-k, 1, k-\alpha+4), & k+1 \leq \alpha \leq k+2, \\ (n-\alpha, n-\alpha+2, \alpha-k, \alpha-k-2, 1), & k+3 \leq \alpha \leq k+4, \\ (n-\alpha, n-\alpha+2, \alpha-k, \alpha-k-2, \alpha-k-4), & k+5 \leq \alpha \leq n-1, \\ (1, 2, k, k-1, k-3), & \alpha = n, \end{cases} \\
 r(a_\alpha b_\alpha | W_1) &= \begin{cases} (0, 2, k, k, k-2), & \alpha = 1, \\ (\alpha, 3-\alpha, k-\alpha+1, k-\alpha+3, \alpha+k-3), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha-2, k-\alpha+1, k-\alpha+3, k-\alpha+5), & 4 \leq \alpha \leq k+1, \\ (n-\alpha+1, \alpha-2, \alpha-k, k-\alpha+3, k-\alpha+5), & k+2 \leq \alpha \leq k+3, \\ (n-\alpha+1, n-\alpha+3, \alpha-k, \alpha-k-2, k-\alpha+5), & k+4 \leq \alpha \leq k+5, \\ (n-\alpha+1, n-\alpha+3, \alpha-k, \alpha-k-2, \alpha-k-4), & k+6 \leq \alpha \leq n, \end{cases} \\
 r(a_\alpha b_{\alpha+1} | W_1) &= \begin{cases} (1, 2, k, k, k-2), & \alpha = 1, \\ (\alpha, \alpha-2, k-\alpha+1, k-\alpha+3, \alpha+k-3), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha-2, k-\alpha+1, k-\alpha+3, k-\alpha+5), & 4 \leq \alpha \leq k-1, \\ (\alpha, \alpha-2, \alpha-k, k-\alpha+3, k-\alpha+5), & k \leq \alpha \leq k+1, \\ (n-\alpha+1, \alpha-2, \alpha-k, \alpha-k-2, k-\alpha+5), & k+2 \leq \alpha \leq k+3, \\ (n-\alpha+1, n-\alpha+3, \alpha-k, \alpha-k-2, \alpha-k-4), & k+4 \leq \alpha \leq n-1, \\ (0, 3, k+1, k-1, k-3), & \alpha = n. \end{cases}
 \end{aligned}$$

From above representation we see that $r(a_\alpha b_\alpha | W_1) = r(a_\alpha b_{\alpha+1} | W_1)$ when $1 \leq \alpha \leq n$ and $\alpha \neq 1, 2, 3, k, k+1, k+2, k+3, k+4, k+5, n$ and no other edges have same representation. If we take odd α , where $1 \leq \alpha \leq n$ and $\alpha \neq 1, 3, k+1, k+3, n-1$ such that $W_E = W_1 \cup \{b_\alpha\}$ then $r(a_\alpha b_\alpha | W_E) \neq r(a_\alpha b_{\alpha+1} | W_E)$ which implies that $\text{edim}(f_{n \times 3}) \leq \lceil \frac{n}{2} \rceil$. So from Lemma 2.1, W_E is an edge metric generator for $f_{n \times 3}$ and $\text{edim}(f_{n \times 3}) = \lceil \frac{n}{2} \rceil$.

Case (IV) When $n \equiv 3 \pmod{4}$. We can write $n = 2k + 1$, $k \geq 3$, k is odd. Let $W_1 = \{b_1, b_3, b_{k+2}, b_{k+4}\}$. Next, we give representation of edges of $f_{n \times 3}$ with respect to W_1 .

$$r(a_\alpha a_{\alpha+1} | W_1) = \begin{cases} (\alpha, 1, k-\alpha+1, \alpha+k-2), & 1 \leq \alpha \leq 2, \\ (\alpha, \alpha-2, k-\alpha+1, k-\alpha+3), & 3 \leq \alpha \leq k-1, \\ (k, \alpha-2, 1, k-\alpha+3), & k \leq \alpha \leq k+1, \\ (n-\alpha, k, \alpha-k-1, 1), & k+2 \leq \alpha \leq k+3, \\ (n-\alpha, n-\alpha+2, \alpha-k-1, \alpha-k-3), & k+4 \leq \alpha \leq n-1, \\ (1, 2, n-k-1, n-k-3), & \alpha = n, \end{cases}$$

$$r(a_\alpha b_\alpha | W_1) = \begin{cases} (0, 2, k+1, n-k-2), & \alpha = 1, \\ (\alpha, 3-\alpha, k-\alpha+2, n-k+\alpha-3), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha-2, k-\alpha+2, k-\alpha+4), & 4 \leq \alpha \leq k+1, \\ (n-k-1, k, 0, 2), & \alpha = k+2, \\ (n-\alpha+1, n-\alpha+3, \alpha-k-1, k-\alpha+4), & k+3 \leq \alpha \leq k+4, \\ (n-\alpha+1, n-\alpha+3, \alpha-k-1, \alpha-k-3), & k+5 \leq \alpha \leq n, \end{cases}$$

$$r(a_\alpha b_{\alpha+1} | W_1) = \begin{cases} (1, 2, k+1, n-k-2), & \alpha = 1, \\ (2, 0, k, n-k-1), & \alpha = 2, \\ (\alpha, \alpha-2, k-\alpha+2, k-\alpha+4), & 3 \leq \alpha \leq k, \\ (n-\alpha+1, \alpha-2, \alpha-k-1, k-\alpha+4), & k+1 \leq \alpha \leq k+2, \\ (n-\alpha+1, n-\alpha+3, \alpha-k-1, \alpha-k-3), & k+3 \leq \alpha \leq n-1, \\ (0, 3, n-k-1, n-k-3), & \alpha = n. \end{cases}$$

From above representation we see that $r(a_\alpha b_\alpha | W_1) = r(a_\alpha b_{\alpha+1} | W_1)$ when $1 \leq \alpha \leq n$ and $\alpha \neq 1, 2, 3, k+1, k+2, k+3, k+4, n$ and no other edges have same representation. If we take odd α , where $1 \leq \alpha \leq n$ and $\alpha \neq 1, 3, k+2, k+4$ such that $W_E = W_1 \cup \{b_\alpha\}$ then $r(a_\alpha b_\alpha | W_E) \neq r(a_\alpha b_{\alpha+1} | W_E)$ which implies that $\text{edim}(f_{n \times 3}) \leq \lceil \frac{n}{2} \rceil$. So from Lemma 2.1, W_E is an edge metric generator for $f_{n \times 3}$ and $\text{edim}(f_{n \times 3}) = \lceil \frac{n}{2} \rceil$. \square

3. Edge metric dimension of D'_n

The prism related graph D'_n has vertex set $V(D'_n) = \{a_\alpha, b_\alpha, c_\alpha, d_\alpha, e_\alpha | 1 \leq \alpha \leq n\}$ and the edge set $E(D'_n) = \{a_\alpha a_{\alpha+1}, a_\alpha b_\alpha, c_\alpha b_\alpha, c_\alpha b_{\alpha+1}, c_\alpha c_{\alpha+1}, c_\alpha d_\alpha, d_\alpha d_{\alpha+1}, d_\alpha e_\alpha | 1 \leq \alpha \leq n\}$ as shown in Figure 2. In this Section, we determine the edge metric of the graph D'_n .

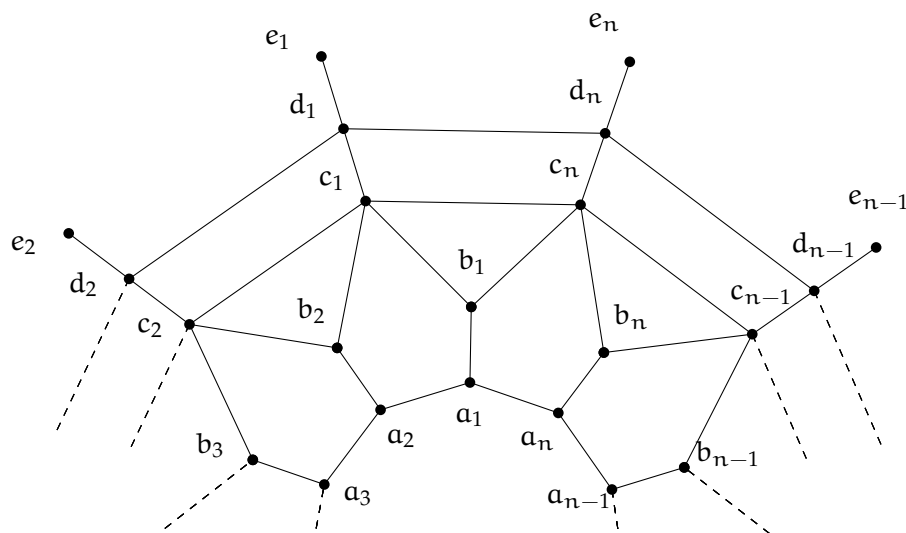


Figure 2: Graph of D'_n .

Theorem 3.1. For the graph D'_n with $n \geq 4$, we have

$$\text{edim}(D'_n) = \begin{cases} 3, & n \text{ is even and } n \neq 4, \\ 4, & n \text{ is odd and } n = 4. \end{cases}$$

Proof. For $n = 4$, we have calculated the edge metric dimension by total enumeration, $\text{edim}(D'_4) = 4$ and its edge metric basis is $\{a_1, a_2, b_3, b_4\}$.

For $n > 4$, we discuss the two cases.

Case (i) When n is even. We can write $n = 2k$, for $k \geq 3$. Let $W_E = \{a_1, e_{k-1}, e_{n-1}\}$ is an edge metric basis for D'_n . We give representations of any edge of $E(D'_n)$ with respect to W_E .

$$\begin{aligned}
 r(a_\alpha a_{\alpha+1} | W_E) &= \begin{cases} (\alpha - 1, k - \alpha + 2, \alpha + 4), & 1 \leq \alpha \leq k - 2, \\ (\alpha - 1, 4, n - \alpha + 2), & k - 1 \leq \alpha \leq k, \\ (n - \alpha, \alpha - k + 4, n - \alpha + 2), & k + 1 \leq \alpha \leq n - 2, \\ (n - \alpha, n - \alpha + k + 2, 4), & n - 1 \leq \alpha \leq n, \end{cases} \\
 r(a_\alpha b_\alpha | W_E) &= \begin{cases} (\alpha - 1, k - \alpha + 2, \alpha + 3), & 1 \leq \alpha \leq k - 1, \\ (k - 1, 3, n - k + 2), & \alpha = k, \\ (n - \alpha + 1, \alpha - k + 3, n - \alpha + 2), & k + 1 \leq \alpha \leq n - 1, \\ (1, k + 2, 3), & \alpha = n, \end{cases} \\
 r(b_\alpha c_\alpha | W_E) &= \begin{cases} (\alpha, k - \alpha + 1, \alpha + 3), & 1 \leq \alpha \leq k - 1, \\ (k, 3, n - k + 1), & \alpha = k, \\ (n - \alpha + 2, \alpha - k + 3, n - \alpha + 1), & k + 1 \leq \alpha \leq n - 1, \\ (2, k + 1, 3), & \alpha = n, \end{cases} \\
 r(c_\alpha b_{\alpha+1} | W_E) &= \begin{cases} (\alpha + 1, k - \alpha + 1, \alpha + 3), & 1 \leq \alpha \leq k - 1, \\ (k + 1, 3, n - k + 1), & \alpha = k, \\ (n - \alpha + 1, \alpha - k + 3, n - \alpha + 1), & k + 1 \leq \alpha \leq n - 1, \\ (1, k + 1, 3), & \alpha = n, \end{cases} \\
 r(c_\alpha c_{\alpha+1} | W_E) &= \begin{cases} (\alpha + 1, k - \alpha, \alpha + 3), & 1 \leq \alpha \leq k - 2, \\ (\alpha + 1, \alpha - k + 3, n - \alpha), & k - 1 \leq \alpha \leq k, \\ (n - \alpha + 1, \alpha - k + 3, n - \alpha), & k + 1 \leq \alpha \leq n - 2, \\ (2, n - \alpha + k, \alpha - n + 3), & n - 1 \leq \alpha \leq n, \end{cases} \\
 r(c_\alpha d_\alpha | W_E) &= \begin{cases} (\alpha + 1, k - \alpha, \alpha + 2), & 1 \leq \alpha \leq k - 1, \\ (k + 1, 2, n - k), & \alpha = k, \\ (n - \alpha + 2, \alpha - k + 2, n - \alpha), & k + 1 \leq \alpha \leq n - 1, \\ (2, k, 2), & \alpha = n, \end{cases} \\
 r(d_\alpha e_\alpha | W_E) &= \begin{cases} (\alpha + 2, k - \alpha, \alpha + 2), & 1 \leq \alpha \leq k - 2, \\ (k + 1, 0, k + 1), & \alpha = k - 1, \\ (k + 2, 2, n - k), & \alpha = k, \\ (n - \alpha + 3, \alpha - k + 2, n - \alpha), & k + 1 \leq \alpha \leq n - 2, \\ (4, n - k + 1, 0), & \alpha = n - 1, \\ (3, k, 2), & \alpha = n, \end{cases} \\
 r(d_\alpha d_{\alpha+1} | W_E) &= \begin{cases} (\alpha + 2, k - \alpha - 1, \alpha + 2), & 1 \leq \alpha \leq k - 2, \\ (\alpha + 2, \alpha - k + 2, n - \alpha - 1), & k - 1 \leq \alpha \leq k, \\ (n - \alpha + 2, \alpha - k + 2, n - \alpha - 1), & k + 1 \leq \alpha \leq n - 2, \\ (3, n - \alpha + k - 1, \alpha - n + 2), & n - 1 \leq \alpha \leq n. \end{cases}
 \end{aligned}$$

We note that there are no vertices having the same edge metric representation implying that $\text{edim}(D_n') \leq 3$. Using Proposition 1.1, $\text{edim}(D_n') \geq 3$, which implies $\text{edim}(D_n') = 3$.

Case (ii) When n is odd. We can write $n = 2k + 1$, for $k \geq 2$. Let $W_E = \{a_1, a_{k+2}, b_2, b_{k+1}\}$ is an edge metric basis for D_n' . We give representations of any edge with respect to W_E .

$$\begin{aligned}
 r(a_\alpha a_{\alpha+1}|W_E) &= \begin{cases} (0, k, 1, k), & \alpha = 1, \\ (\alpha - 1, k - \alpha + 1, \alpha - 1, k - \alpha + 1), & 2 \leq \alpha \leq k, \\ (n - k - 1, 0, k, 1), & \alpha = k + 1, \\ (n - \alpha, \alpha - k - 2, n - \alpha + 2, \alpha - k), & k + 2 \leq \alpha \leq n, \end{cases} \\
 r(a_\alpha b_\alpha|W_E) &= \begin{cases} (0, k, 2, k + 1), & \alpha = 1, \\ (1, k, 0, k), & \alpha = 2, \\ (\alpha - 1, k - \alpha + 2, \alpha - 1, k - \alpha + 2), & 3 \leq \alpha \leq k, \\ (k, 1, k, 0), & \alpha = k + 1, \\ (n - k - 1, 0, k + 1, 2), & \alpha = k + 2, \\ (n - \alpha + 1, \alpha - k - 2, n - \alpha + 3, \alpha - k), & k + 3 \leq \alpha \leq n, \end{cases} \\
 r(b_\alpha c_\alpha|W_E) &= \begin{cases} (1, k + 1, 1, k), & \alpha = 1, \\ (2, k + 1, 0, k - 1), & \alpha = 2, \\ (\alpha, k - \alpha + 3, \alpha - 1, k - \alpha + 1), & 3 \leq \alpha \leq k + 1, \\ (n - \alpha + 2, \alpha - k - 1, n - \alpha + 2, \alpha - k), & k + 2 \leq \alpha \leq n, \end{cases} \\
 r(c_\alpha b_{\alpha+1}|W_E) &= \begin{cases} (\alpha + 1, k - \alpha + 2, \alpha - 1, k - \alpha + 1), & 1 \leq \alpha \leq k - 1, \\ (k + 1, 2, k - 1, 0), & \alpha = k, \\ (k + 2, 1, k, 1), & \alpha = k + 1, \\ (n - \alpha + 1, \alpha - k, n - \alpha + 2, \alpha - k), & k + 2 \leq \alpha \leq n, \end{cases} \\
 r(c_\alpha c_{\alpha+1}|W_E) &= \begin{cases} (2, k + 1, 1, k - 1), & \alpha = 1, \\ (\alpha + 1, k - \alpha + 2, \alpha - 1, k - \alpha), & 2 \leq \alpha \leq k - 1, \\ (k + 1, 2, \alpha - 1, 1), & k \leq \alpha \leq k + 1, \\ (n - \alpha + 1, \alpha - k, n - \alpha + 1, \alpha - k), & k + 2 \leq \alpha \leq n - 1, \\ (2, k + 1, 1, k), & \alpha = n, \end{cases} \\
 r(c_\alpha d_\alpha|W_E) &= \begin{cases} (2, k + 2, 1, k), & \alpha = 1, \\ (\alpha + 1, k - \alpha + 3, \alpha - 1, k - \alpha + 1), & 2 \leq \alpha \leq k, \\ (k + 2, 2, k, 1), & \alpha = k + 1, \\ (n - \alpha + 2, \alpha - k, n - \alpha + 2, \alpha - k), & k + 2 \leq \alpha \leq n, \end{cases} \\
 r(d_\alpha d_{\alpha+1}|W_E) &= \begin{cases} (3, k + 2, 2, k), & \alpha = 1, \\ (\alpha + 2, k - \alpha + 3, \alpha, k - \alpha + 1), & 2 \leq \alpha \leq k - 1, \\ (k + 2, 3, \alpha, 2), & k \leq \alpha \leq k + 1, \\ (n - \alpha + 2, \alpha - k + 1, n - \alpha + 2, \alpha - k + 1), & k + 2 \leq \alpha \leq n - 1, \\ (3, k + 2, 2, k + 1), & \alpha = n, \end{cases} \\
 r(d_\alpha e_\alpha|W_E) &= \begin{cases} (3, k + 3, 2, k + 1), & \alpha = 1, \\ (\alpha + 2, k - \alpha + 4, \alpha, k - \alpha + 2), & 2 \leq \alpha \leq k, \\ (k + 3, 3, k + 1, 2), & \alpha = k + 1, \\ (n - \alpha + 3, \alpha - k + 1, n - \alpha + 3, \alpha - k + 1), & k + 2 \leq \alpha \leq n. \end{cases}
 \end{aligned}$$

We note that there are no vertices having the same edge metric representation implying that $\text{edim}(D_n') \leq 4$.

Suppose on contrary that $\text{edim}(D_n') = 3$, then the Table 1 shows all order pairs of edges (e, f) for which $r(e|W_E) = r(f|W_E)$.

Table 1: (e, f) for which $r(e|W_E) = r(f|W_E)$.

W_E	(e, f)
$\{a_1, a_j, a_l\}$ where $2 \leq j \leq k+1$ and $3 \leq l \leq n$	For $3 \leq l \leq k+1$, $r(a_1 a_n W_E) = r(a_1 b_1 W_E)$. And for $k+2 \leq l \leq n$, $r(b_1 c_1 W_E) = r(b_1 c_n W_E)$.
$\{b_1, b_j, b_l\}$ where $2 \leq j \leq k+1$ and $3 \leq l \leq n$	For $3 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_n W_E) = r(d_n e_n W_E)$.
$\{c_1, c_j, c_l\}$ where $2 \leq j \leq k+1$ and $3 \leq l \leq n$	For $3 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_2 W_E) = r(d_1 e_1 W_E)$.
$\{d_1, d_j, d_l\}$ where $2 \leq j \leq k+1$ and $3 \leq l \leq n$	This case is same to case (3) due to symmetry of the graph.
$\{e_1, e_j, e_l\}$ where $2 \leq j \leq k+1$ and $3 \leq l \leq n$	For $3 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(c_1 d_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_2 W_E) = r(c_1 d_1 W_E)$.
$\{a_1, a_j, b_l\}$ where $2 \leq j \leq k+1$ and $1 \leq l \leq n$	For $1 \leq l \leq k+1$, $r(e_k d_k W_E) = r(d_k d_{k+1} W_E)$. For $k+2 \leq l \leq n-1$, $r(b_2 c_2 W_E) = r(c_1 b_2 W_E)$. And for $l = n$, $r(c_{k+1} d_{k+1} W_E) = r(d_k d_{k+1} W_E)$.
$\{a_1, a_j, c_l\}$ where $2 \leq j \leq k+1$ and $1 \leq l \leq n$	For $1 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_2 W_E) = r(d_1 e_1 W_E)$.
$\{a_1, a_j, e_l\}$ where $2 \leq j \leq k+1$ and $1 \leq l \leq n$	For $1 \leq l \leq k$, $r(d_{n-1} d_n W_E) = r(d_n e_n W_E)$. For $l = k+1$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_2 W_E) = r(d_1 e_1 W_E)$.
$\{a_1, b_j, b_l\}$ where $1 \leq j \leq k+1$ and $2 \leq l \leq n$	For $2 \leq l \leq k$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. For $k+2 \leq l \leq n-1$, $r(c_1 b_2 W_E) = r(b_2 c_2 W_E)$. And for $l = n$, $r(d_1 d_2 W_E) = r(d_1 e_1 W_E)$.
$\{a_1, b_j, c_l\}$ where $1 \leq j \leq k+1$ and $1 \leq l \leq n$	For $1 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_2 W_E) = r(d_1 e_1 W_E)$.
$\{a_1, b_j, e_l\}$ where $1 \leq j \leq k+1$ and $1 \leq l \leq n$	For $1 \leq l \leq k$, $r(d_{n-1} d_n W_E) = r(d_n e_n W_E)$. For $l = k+1$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_2 W_E) = r(d_1 e_1 W_E)$.
$\{a_1, c_j, c_l\}$ where $1 \leq j \leq k+1$ and $2 \leq l \leq n$	For $2 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_2 W_E) = r(d_1 e_1 W_E)$.
$\{a_1, c_j, d_l\}$ where $1 \leq j \leq k+1$ and $1 \leq l \leq n$	For $1 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(d_1 e_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_1 d_2 W_E) = r(d_1 e_1 W_E)$.
$\{a_1, e_j, e_l\}$ where $1 \leq j \leq k+1$ and $2 \leq l \leq n$	For $2 \leq l \leq k$, $r(d_n d_{n-1} W_E) = r(d_n e_n W_E)$. For $l = k+1$, $r(d_{n-1} d_n W_E) = r(c_2 d_2 W_E)$. And for $k+2 \leq l \leq n$, $r(d_{k+1} c_{k+1} W_E) = r(c_{k+1} b_{k+2} W_E)$.
$\{b_1, e_j, e_l\}$ where $1 \leq j \leq k+1$ and $2 \leq l \leq n$	For $2 \leq l \leq k$, $r(d_n d_{n-1} W_E) = r(d_n e_n W_E)$. And for $k+1 \leq l \leq n$, $r(d_2 e_2 W_E) = r(c_2 d_2 W_E)$.
$\{d_1, d_j, e_l\}$ where $2 \leq j \leq k+1$ and $1 \leq l \leq n$	For $1 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(c_1 d_1 W_E)$. And for $k+2 \leq l \leq n$, $r(d_2 e_2 W_E) = r(c_2 d_2 W_E)$.
$\{d_1, e_j, e_l\}$ where $1 \leq j \leq k+1$ and $2 \leq l \leq n$	For $2 \leq l \leq k+1$, $r(d_1 d_n W_E) = r(c_1 d_1 W_E)$. For $k+2 \leq l \leq n$, $r(d_2 e_2 W_E) = r(c_2 d_2 W_E)$.

Hence, Table 1 follows that there is no resolving set with three vertices for $V(D_n')$ implying that $\text{edim}(D_n') = 4$. Which completes the proof. \square

4. Edge metric dimension of D_n^t

The prism related graph D_n^t has vertex set $V(D_n^t) = \{a_\alpha, b_\alpha, c_\alpha, d_\alpha : 1 \leq \alpha \leq n\}$ and the edge set $E(D_n^t) = \{a_\alpha a_{\alpha+1}, a_\alpha b_\alpha, b_\alpha b_{\alpha+1}, b_\alpha c_\alpha, b_\alpha c_{\alpha+1}, c_\alpha d_\alpha : 1 \leq \alpha \leq n\}$ as shown in Figure 3. In this Section, we determine the edge metric dimension of the graph D_n^t .

Lemma 4.1. *Let $Y = \{d_1, d_2, \dots, d_n\}$ be a subset of $V(D_n^t)$. Then any arbitrary edge metric generator W_E of D_n^t contains at least $\lceil \frac{n}{2} \rceil$ vertices.*

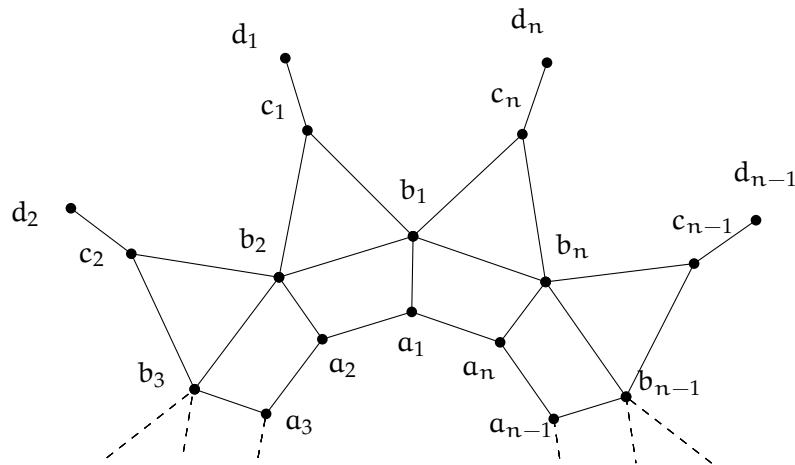


Figure 3: Graph of D_n^t .

Proof. Suppose that W_E contains at most $\lceil \frac{n}{2} \rceil - 1$ vertices of Y for a contradiction. Without loss of generality we assume that $d_\alpha, d_{\alpha+1} \notin W_E$, then we have $(b_\alpha c_\alpha | W_E) = (b_\alpha c_{\alpha+1} | W_E)$, a contradiction. \square

Remark 4.2. Let W_E be any edge metric basis of D_n^t . We note that W_E contains all odd vertices (vertex indices are odd) of Y for odd n , while W_E contains either all odd vertices or even vertices (vertex indices are even) of Y for even n .

Lemma 4.3. For $n \geq 5$, we have $\text{edim}(D_n^t) \geq \lceil \frac{n}{2} \rceil + 1$

Proof. We assume for a contradiction that the cardinality of subset W_E is equal to $\lceil \frac{n}{2} \rceil$ by Lemma 4.1. Using Remark 4.2, we take $W_E = \{d_\alpha \in W_E | \text{Vertices indices } \alpha \text{ is odd}\}$ such that $|W_E| = \lceil \frac{n}{2} \rceil$. We have $(a_\alpha b_\alpha | W_E) = (b_\alpha c_\alpha | W_E)$ for even α ($1 \leq \alpha \leq n$), a contradiction. So, $\text{edim}(D_n^t) \geq \lceil \frac{n}{2} \rceil + 1$. \square

Theorem 4.4. For the graph D_n^t with $n \geq 3$, we have

$$\text{edim}(D_n^t) = \begin{cases} 4, & n = 3, 4, \\ \lceil \frac{n}{2} \rceil + 1, & \text{otherwise.} \end{cases}$$

Proof. For $n = 3, 4$ we have calculated the edge metric dimension by total enumeration that is 4 and its edge metric basis is $\{c_1, c_2, c_3, c_4\}$ and $\{a_1, b_3, c_1, c_2\}$, respectively.

For $n \geq 5$, we discuss the following four cases.

Let $W_E = \{a_1, d_1, d_3, d_5, \dots, d_{n-1}\}$. We will show that W_E is an edge metric basis of D_n^t in Case(I) and Case(II), respectively.

Case (I) When $n \equiv 0 \pmod{4}$. We can write $n = 2k$, $k \geq 4$, and k is even. Let $W_1 = \{a_1, d_1, d_3, d_{k+1}, d_{k+3}\}$. Next, we give representation of edges of D_n^t with respect to W_1 .

$$r(a_\alpha a_{\alpha+1} | W_1) = \begin{cases} (\alpha - 1, \alpha + 2, 3, k - \alpha + 2, \alpha + k), & 1 \leq \alpha \leq 2, \\ (\alpha - 1, \alpha + 2, \alpha, k - \alpha + 2, k - \alpha + 4), & 3 \leq \alpha \leq k - 1, \\ (k - 1, k + 2, k, 3, 4), & \alpha = k, \\ (n - \alpha, n - \alpha + 2, \alpha, \alpha - k + 2, 3), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha, n - \alpha + 2, n - \alpha + 4, \alpha - k + 2, \alpha - k), & k + 3 \leq \alpha \leq n - 1, \\ (0, 3, 4, k + 2, k), & \alpha = n, \end{cases}$$

$$r(a_\alpha b_\alpha | W_1) = \begin{cases} (\alpha - 1, \alpha + 1, 4 - \alpha, k - \alpha + 2, \alpha + k - 1), & 1 \leq \alpha \leq 2, \\ (\alpha - 1, \alpha + 1, \alpha - 1, k - \alpha + 2, k - \alpha + 4), & 3 \leq \alpha \leq k, \\ (n - \alpha + 1, n - \alpha + 2, \alpha - 1, \alpha - k + 1, k - \alpha + 4), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha + 1, n - \alpha + 2, n - \alpha + 4, \alpha - k + 1, \alpha - k - 1), & k + 3 \leq \alpha \leq n, \end{cases}$$

$$\begin{aligned}
 r(b_\alpha b_{\alpha+1}|W_1) &= \begin{cases} (\alpha, \alpha + 1, 2, k - \alpha + 1, \alpha + k - 1), & 1 \leq \alpha \leq 2, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 1, k - \alpha + 3), & 3 \leq \alpha \leq k - 1, \\ (k, k + 1, k - 1, 2, 3), & \alpha = k, \\ (n - \alpha + 1, n - \alpha + 1, \alpha - 1, \alpha - k + 1, 2), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha + 1, n - \alpha + 1, n - \alpha + 3, \alpha - k + 1, \alpha - k - 1), & k + 3 \leq \alpha \leq n - 1, \\ (1, 2, 3, k + 1, k - 1), & \alpha = n, \end{cases} \\
 r(b_\alpha c_\alpha|W_1) &= \begin{cases} (1, 1, 3, k + 1, k), & \alpha = 1, \\ (\alpha, \alpha + 1, 4 - \alpha, k - \alpha + 2, k + 1), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 2, k - \alpha + 4), & 4 \leq \alpha \leq k, \\ (n - k + 1, n - k + 1, k, 1, 3), & \alpha = k + 1, \\ (n - k, n - k, k + 1, 3, 2), & \alpha = k + 2, \\ (n - k - 1, n - k - 1, n - k + 1, 4, 1), & \alpha = k + 3, \\ (n - \alpha + 2, n - \alpha + 2, n - \alpha + 4, \alpha - k + 1, \alpha - k - 1), & k + 4 \leq \alpha \leq n, \end{cases} \\
 r(b_\alpha c_{\alpha+1}|W_1) &= \begin{cases} (1, 2, 3, k + 1, k), & \alpha = 1, \\ (2, 3, 1, k, k + 1), & \alpha = 2, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 2, k - \alpha + 4), & 3 \leq \alpha \leq k - 1, \\ (k, k + 1, k - 1, 1, 4), & \alpha = k, \\ (n - k + 1, n - k + 1, k, 2, 3), & \alpha = k + 1, \\ (n - k, n - k, k + 1, 3, 1), & \alpha = k + 2, \\ (n - \alpha + 2, n - \alpha + 2, n - \alpha + 4, \alpha - k + 1, \alpha - k - 1), & k + 3 \leq \alpha \leq n - 1, \\ (2, 1, 4, k + 1, k - 1), & \alpha = n, \end{cases} \\
 r(c_\alpha d_\alpha|W_1) &= \begin{cases} (2, 0, 4, k + 2, k), & \alpha = 1, \\ (2, 3, 3, k + 1, k + 1), & \alpha = 2, \\ (3, 4, 0, k, k + 2), & \alpha = 3, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 3, k - \alpha + 5), & 4 \leq \alpha \leq k, \\ (k + 1, k + 2, k, 0, 4), & \alpha = k + 1, \\ (n - k + 1, n - k + 1, k + 1, 3, 3), & \alpha = k + 2, \\ (n - k, n - k, n - k + 2, 4, 0), & \alpha = k + 3, \\ (n - \alpha + 3, n - \alpha + 3, n - \alpha + 5, \alpha - k + 1, \alpha - k - 1), & k + 4 \leq \alpha \leq n. \end{cases}
 \end{aligned}$$

From above representation we see that $r(b_\alpha c_\alpha|W_1) = r(b_\alpha c_{\alpha+1}|W_1)$ when $1 \leq \alpha \leq n$ and $\alpha \neq 1, 2, 3, k, k + 1, k + 2, k + 3, n$ and no other edges have same representation. If we take odd α , where $1 \leq \alpha \leq n$ and $\alpha \neq 1, 3, k + 1, n - 1$ such that $W_E = W_1 \cup \{d_\alpha\}$ then $r(b_\alpha c_\alpha|W_E) \neq r(b_\alpha c_{\alpha+1}|W_E)$ which implies that $\text{edim}(D_n^t) \leq \lceil \frac{n}{2} \rceil + 1$. So from Lemma 4.3 W_E is an edge metric basis for D_n^t and $\text{edim}(D_n^t) = \lceil \frac{n}{2} \rceil + 1$.

Case (II) When $n \equiv 2 \pmod{4}$. We can write $n = 2k$, $k \geq 3$, and k is odd. Let $W_1 = \{a_1, d_1, d_3, d_{k+2}\}$. Next, we give representation of edges of D_n^t with respect to W_1 .

$$\begin{aligned}
 r(a_\alpha a_{\alpha+1}|W_1) &= \begin{cases} (\alpha - 1, \alpha + 2, 3, k - \alpha + 3), & 1 \leq \alpha \leq 2, \\ (\alpha - 1, \alpha + 2, \alpha, k - \alpha + 3), & 3 \leq \alpha \leq k, \\ (k - 1, k + 1, k + 1, 3), & \alpha = k + 1, \\ (n - \alpha, n - \alpha + 2, n - \alpha + 4, \alpha - k + 1), & k + 2 \leq \alpha \leq n - 1, \\ (0, 3, 4, k + 1), & \alpha = n, \end{cases} \\
 r(a_\alpha b_\alpha|W_1) &= \begin{cases} (\alpha - 1, \alpha + 1, 4 - \alpha, k + 1), & 1 \leq \alpha \leq 2, \\ (\alpha - 1, \alpha + 1, \alpha - 1, k - \alpha + 3), & 3 \leq \alpha \leq k, \\ (n - \alpha + 1, n - \alpha + 2, \alpha - 1, 2), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha + 1, n - \alpha + 2, n - \alpha + 4, \alpha - k), & k + 3 \leq \alpha \leq n, \end{cases}
 \end{aligned}$$

$$\begin{aligned}
 r(b_\alpha b_{\alpha+1}|W_1) &= \begin{cases} (\alpha, \alpha + 1, 2, k - \alpha + 2), & 1 \leq \alpha \leq 2, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 2), & 3 \leq \alpha \leq k, \\ (k, k, k, 2), & \alpha = k + 1, \\ (n - \alpha + 1, n - \alpha + 1, n - \alpha + 3, \alpha - k), & k + 2 \leq \alpha \leq n - 1, \\ (1, 2, 3, k), & \alpha = n, \end{cases} \\
 r(b_\alpha c_\alpha|W_1) &= \begin{cases} (1, 1, 3, k + 1), & \alpha = 1, \\ (\alpha, \alpha + 1, 4 - \alpha, k - \alpha + 3), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 3), & 4 \leq \alpha \leq k, \\ (n - \alpha + 2, n - \alpha + 2, \alpha - 1, k - \alpha + 3), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha + 2, n - \alpha + 2, n - \alpha + 4, \alpha - k), & k + 3 \leq \alpha \leq n, \end{cases} \\
 r(b_\alpha c_{\alpha+1}|W_1) &= \begin{cases} (1, 2, 3, k + 1), & \alpha = 1, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 3), & 2 \leq \alpha \leq k, \\ (n - \alpha + 2, n - \alpha + 2, \alpha - 1, \alpha - k), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha + 2, n - \alpha + 2, n - \alpha + 4, \alpha - k), & k + 3 \leq \alpha \leq n - 1, \\ (2, 1, 4, k), & \alpha = n, \end{cases} \\
 r(c_\alpha d_\alpha|W_1) &= \begin{cases} (2, 0, 4, k + 1), & \alpha = 1, \\ (2, 3, 3, k + 2), & \alpha = 2, \\ (3, 4, 0, k + 1), & \alpha = 3, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 4), & 4 \leq \alpha \leq k + 1, \\ (n - k + 1, n - k + 1, k + 1, 0), & \alpha = k + 2, \\ (n - \alpha + 3, n - \alpha + 3, n - \alpha + 5, \alpha - k), & k + 3 \leq \alpha \leq n. \end{cases}
 \end{aligned}$$

From above representation we see that $r(b_\alpha c_\alpha|W_1) = r(b_\alpha c_{\alpha+1}|W_1)$ when $1 \leq \alpha \leq n$ and $\alpha \neq 1, 2, 3, k + 1, k + 2, n$ and no other edges have same representation. If we take odd α , where $1 \leq \alpha \leq n$ and $\alpha \neq 1, 3, k + 2$ such that $W_E = W_1 \cup \{d_\alpha\}$ then $r(b_\alpha c_\alpha|W_E) \neq r(b_\alpha c_{\alpha+1}|W_E)$ which implies that $\text{edim}(D_n^t) \leq \lceil \frac{n}{2} \rceil + 1$. So from Lemma 4.3 W_E is an edge metric basis for D_n^t and $\text{edim}(D_n^t) = \lceil \frac{n}{2} \rceil + 1$.

Let $W_E = \{a_1, d_1, d_3, d_5, \dots, d_{n-2}, d_n\}$. We will show that W_E is an edge metric basis of D_n^t in Cases (III) and (IV), respectively.

Case (III) When $n \equiv 1 \pmod{4}$. We can write $n = 2k + 1, k \geq 2$, and k is even. Let $W_1 = \{a_1, d_1, d_3, d_{k+3}\}$. Next, we give representation of edges of D_n^t with respect to W_1 .

$$\begin{aligned}
 r(a_\alpha a_{\alpha+1}|W_1) &= \begin{cases} (\alpha - 1, \alpha + 2, 3, k + 2), & 1 \leq \alpha \leq 2, \\ (\alpha - 1, \alpha + 2, \alpha, k - \alpha + 4), & 3 \leq \alpha \leq k, \\ (n - \alpha, n - \alpha + 2, \alpha, 3), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha, n - \alpha + 2, n - \alpha + 4, \alpha - k), & k + 3 \leq \alpha \leq n - 1, \\ (0, 3, 4, k + 1), & \alpha = n, \end{cases} \\
 r(a_\alpha b_\alpha|W_1) &= \begin{cases} (\alpha - 1, \alpha + 1, 4 - \alpha, k + \alpha), & 1 \leq \alpha \leq 2, \\ (\alpha - 1, \alpha + 1, \alpha - 1, k - \alpha + 4), & 3 \leq \alpha \leq k + 1, \\ (k, k + 1, k + 1, 2), & \alpha = k + 2, \\ (n - \alpha + 1, n - \alpha + 2, n - \alpha + 4, \alpha - k - 1), & k + 3 \leq \alpha \leq n, \end{cases} \\
 r(b_\alpha b_{\alpha+1}|W_1) &= \begin{cases} (\alpha, \alpha + 1, 2, k + 1), & 1 \leq \alpha \leq 2, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 3), & 3 \leq \alpha \leq k, \\ (n - \alpha + 1, n - \alpha + 1, \alpha - 1, 2), & k + 1 \leq \alpha \leq k + 2, \\ (n - \alpha + 1, n - \alpha + 1, n - \alpha + 3, \alpha - k - 1), & k + 3 \leq \alpha \leq n - 1, \\ (1, 2, 3, k), & \alpha = n, \end{cases}
 \end{aligned}$$

$$\begin{aligned}
 r(b_\alpha c_\alpha | W_1) &= \begin{cases} (1, 1, 3, k+1), & \alpha = 1, \\ (\alpha, \alpha + 1, 4 - \alpha, k - \alpha + 4), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 4), & 4 \leq \alpha \leq k+1, \\ (n - \alpha + 2, n - \alpha + 2, \alpha - 1, k - \alpha + 4), & k+2 \leq \alpha \leq k+3, \\ (n - \alpha + 2, n - \alpha + 2, n - \alpha + 4, \alpha - k - 1), & k+4 \leq \alpha \leq n, \end{cases} \\
 r(b_\alpha c_{\alpha+1} | W_1) &= \begin{cases} (1, 2, 3, k+1), & \alpha = 1, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 4), & 2 \leq \alpha \leq k+1, \\ (n - \alpha + 2, n - \alpha + 2, \alpha - 1, \alpha - k - 1), & k+2 \leq \alpha \leq k+3, \\ (n - \alpha + 2, n - \alpha + 2, n - \alpha + 4, \alpha - k - 1), & k+4 \leq \alpha \leq n-1, \\ (2, 1, 4, k), & \alpha = n, \end{cases} \\
 r(c_\alpha d_\alpha | W_1) &= \begin{cases} (2, 0, 4, k+1), & \alpha = 1, \\ (2, 3, 3, k+2), & \alpha = 2, \\ (3, 4, 0, k+2), & \alpha = 3, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 5), & 4 \leq \alpha \leq k+1, \\ (k+2, n - k + 1, k + 1, 3), & \alpha = k+2, \\ (n - k, n - k, k + 2, 0), & \alpha = k+3, \\ (n - \alpha + 3, n - \alpha + 3, n - \alpha + 5, \alpha - k - 1), & k+4 \leq \alpha \leq n. \end{cases}
 \end{aligned}$$

From above representation we see that $r(b_\alpha c_\alpha | W_1) = r(b_\alpha c_{\alpha+1} | W_1)$ when $1 \leq \alpha \leq n$ and $\alpha \neq 1, 2, 3, k+2, k+3, n$ and no other edges have same representation. If we take odd α , where $1 \leq \alpha \leq n$ and $\alpha \neq 1, 3, k+3$ such that $W_E = W_1 \cup \{d_\alpha\}$ then $r(b_\alpha c_\alpha | W_E) \neq r(b_\alpha c_{\alpha+1} | W_E)$ which implies that $\text{edim}(D_n^t) \leq \lceil \frac{n}{2} \rceil + 1$. So from Lemma 4.3 W_E is an edge metric basis for D_n^t and $\text{edim}(D_n^t) = \lceil \frac{n}{2} \rceil + 1$.

Case (IV) When $n \equiv 3 \pmod{4}$. We can write $n = 2k + 1$, $k \geq 3$, and k is odd. Let $W_1 = \{a_1, d_1, d_3, d_{k+2}, d_{k+4}\}$. Next, we give representation of edges of D_n^t with respect to W_1 .

$$\begin{aligned}
 r(a_\alpha a_{\alpha+1} | W_1) &= \begin{cases} (\alpha - 1, \alpha + 2, 3, k - \alpha + 3, \alpha + k), & 1 \leq \alpha \leq 2, \\ (\alpha - 1, \alpha + 2, \alpha, k - \alpha + 3, k - \alpha + 5), & 3 \leq \alpha \leq k, \\ (n - \alpha, n - \alpha + 2, \alpha, 3, k - \alpha + 5), & k+1 \leq \alpha \leq k+2, \\ (n - k - 3, n - k - 1, n - k + 1, 4, 3), & \alpha = k+3, \\ (n - \alpha, n - \alpha + 2, n - \alpha + 4, \alpha - k + 1, \alpha - k - 1), & k+4 \leq \alpha \leq n-1, \\ (0, 3, 4, n - k + 1, n - k - 1), & \alpha = n, \end{cases} \\
 r(a_\alpha b_\alpha | W_1) &= \begin{cases} (\alpha - 1, \alpha + 1, 4 - \alpha, k - \alpha + 3, k + \alpha - 1), & 1 \leq \alpha \leq 2, \\ (\alpha - 1, \alpha + 1, \alpha - 1, k - \alpha + 3, k - \alpha + 5), & 3 \leq \alpha \leq k+1, \\ (n - \alpha + 1, n - \alpha + 2, \alpha - 1, \alpha - k, k - \alpha + 5), & k+2 \leq \alpha \leq k+3, \\ (n - \alpha + 1, n - \alpha + 2, n - \alpha + 4, \alpha - k, \alpha - k - 2), & k+4 \leq \alpha \leq n, \end{cases} \\
 r(b_\alpha b_{\alpha+1} | W_1) &= \begin{cases} (\alpha, \alpha + 1, 2, k - \alpha + 2, \alpha + k - 1), & 1 \leq \alpha \leq 2, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 2, k - \alpha + 4), & 3 \leq \alpha \leq k, \\ (k + 1, n - k, k, 2, 3), & \alpha = k+1, \\ (n - \alpha + 1, n - \alpha + 1, k + 1, \alpha - k, 2), & k+2 \leq \alpha \leq k+3, \\ (n - \alpha + 1, n - \alpha + 1, n - \alpha + 3, \alpha - k, \alpha - k - 2), & k+4 \leq \alpha \leq n-1, \\ (1, 2, 3, n - k, n - k - 2), & \alpha = n, \end{cases} \\
 r(b_\alpha c_\alpha | W_1) &= \begin{cases} (1, 1, 3, k+2, k), & \alpha = 1, \\ (\alpha, \alpha + 1, 4 - \alpha, k - \alpha + 3, \alpha + k - 1), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha + 1, \alpha - 1, k - \alpha + 3, k - \alpha + 5), & 4 \leq \alpha \leq k+1, \\ (n - k, n - k, k + 1, 1, 3), & \alpha = k+2, \\ (n - \alpha + 2, n - \alpha + 2, n - \alpha + 4, \alpha - k, k - \alpha + 5), & k+3 \leq \alpha \leq k+4, \\ (n - \alpha + 2, n - \alpha + 2, n - \alpha + 4, \alpha - k, \alpha - k - 2), & k+5 \leq \alpha \leq n, \end{cases}
 \end{aligned}$$

$$r(b_{\alpha}c_{\alpha+1}|W_1) = \begin{cases} (1, 2, 3, k+2, k), & \alpha = 1, \\ (\alpha, \alpha+1, \alpha-1, k-\alpha+3, k+\alpha-1), & 2 \leq \alpha \leq 3, \\ (\alpha, \alpha+1, \alpha-1, k-\alpha+3, k-\alpha+5), & 4 \leq \alpha \leq k, \\ (k+1, n-\alpha+2, \alpha-1, \alpha-k, k-\alpha+5), & k+1 \leq \alpha \leq k+2, \\ (n-\alpha+2, n-\alpha+2, n-\alpha+4, \alpha-k, \alpha-k-2), & k+3 \leq \alpha \leq n-1, \\ (2, 1, 4, k+1, k-1), & \alpha = n, \end{cases}$$

$$r(c_{\alpha}d_{\alpha}|W_1) = \begin{cases} (2, 0, 4, k+2, k), & \alpha = 1, \\ (2, 3, 3, k+2, k+1), & \alpha = 2, \\ (3, 4, 0, k+1, k+2), & \alpha = 3, \\ (\alpha, \alpha+1, \alpha-1, k-\alpha+4, k-\alpha+6), & 4 \leq \alpha \leq k+1, \\ (n-k+1, n-k+1, k+1, 0, 4), & \alpha = k+2, \\ (n-k, n-k, k+2, 3, 3), & \alpha = k+3, \\ (n-k-1, n-k-1, n-k+1, 4, 0), & \alpha = k+4, \\ (n-\alpha+3, n-\alpha+3, n-\alpha+5, \alpha-k, \alpha-k-2), & k+5 \leq \alpha \leq n. \end{cases}$$

From above representation we see that $r(b_{\alpha}c_{\alpha}|W_1) = r(b_{\alpha}c_{\alpha+1}|W_1)$ when $1 \leq \alpha \leq n$ and $\alpha \neq 1, 2, 3, k+1, k+2, k+3, k+4, n$ and no other edges have same representation. If we take odd α , where $1 \leq \alpha \leq n$ and $\alpha \neq 1, 3, k+2, k+4$ such that $W_E = W_1 \cup \{d_{\alpha}\}$ then $r(b_{\alpha}c_{\alpha}|W_E) \neq r(b_{\alpha}c_{\alpha+1}|W_E)$ which implies that $\text{edim}(D_n^t) \leq \lceil \frac{n}{2} \rceil + 1$. So from Lemma 4.3 W_E is an edge metric basis for D_n^t and $\text{edim}(D_n^t) = \lceil \frac{n}{2} \rceil + 1$. \square

5. Conclusion

In this paper, we have computed the exact value of the edge metric dimension of flower graph $f_{n \times 3}$, the prism related graphs D_n' and D_n^t . It has been observed that the edge metric dimension of these graphs is greater than the metric dimension and we concluded that the edge metric dimension of D_n' is constant while $f_{n \times 3}$ and D_n^t have unbounded edge metric dimensions.

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